

An introduction to polynomial interpolation

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- 1 Introduction
- 2 Interpolation on an arbitrary grid
- 3 Expansions onto orthogonal polynomials
- 4 Convergence of the spectral expansions
- 5 References

Outline

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Introduction

Basic idea: approximate functions $\mathbb{R} \rightarrow \mathbb{R}$ by **polynomials**

Polynomials are the only functions that a computer can evaluate exactly.

Two types of numerical methods based on polynomial approximations:

- **spectral methods:** high order polynomials on a single domain (or a few domains)
- **finite elements:** low order polynomials on many domains

Introduction

Basic idea: approximate functions $\mathbb{R} \rightarrow \mathbb{R}$ by **polynomials**

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- **spectral methods**: high order polynomials on a single domain (or a few domains)
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Framework of this lecture

We consider real-valued functions on the compact interval $[-1, 1]$:

$$f : [-1, 1] \longrightarrow \mathbb{R}$$

We denote

- by \mathbb{P} the set all real-valued polynomials on $[-1, 1]$:

$$\forall p \in \mathbb{P}, \forall x \in [-1, 1], p(x) = \sum_{i=0}^n a_i x^i$$

- by \mathbb{P}_N (where N is a positive integer), the subset of polynomials of degree at most N .

Is it a good idea to approximate functions by polynomials ?

For **continuous functions**, the answer is **yes**:

Theorem (Weierstrass, 1885)

\mathbb{P} is a dense subspace of the space $C^0([-1, 1])$ of all continuous functions on $[-1, 1]$, equipped with the uniform norm $\|.\|_\infty$.^a

^aThis is a particular case of the Stone-Weierstrass theorem

The **uniform norm** or **maximum norm** is defined by $\|f\|_\infty = \max_{x \in [-1, 1]} |f(x)|$

Other phrasings:

For any continuous function on $[-1, 1]$, f , and any $\epsilon > 0$, there exists a polynomial $p \in \mathbb{P}$ such that $\|f - p\|_\infty < \epsilon$.

For any continuous function on $[-1, 1]$, f , there exists a sequence of polynomials $(p_n)_{n \in \mathbb{N}}$ which converges uniformly towards f : $\lim_{n \rightarrow \infty} \|f - p_n\|_\infty = 0$.

Best approximation polynomial

For a given continuous function: $f \in C^0([-1, 1])$, a **best approximation polynomial of degree N** is a polynomial $p_N^*(f) \in \mathbb{P}_N$ such that

$$\|f - p_N^*(f)\|_\infty = \min \{\|f - p\|_\infty, p \in \mathbb{P}_N\}$$

Chebyshev's alternant theorem (or equioscillation theorem)

For any $f \in C^0([-1, 1])$ and $N \geq 0$, the best approximation polynomial $p_N^*(f)$ exists and is unique. Moreover, there exists $N + 2$ points x_0, x_1, \dots, x_{N+1} in $[-1, 1]$ such that

$$f(x_i) - p_N^*(f)(x_i) = (-1)^i \|f - p_N^*(f)\|_\infty, \quad 0 \leq i \leq N + 1$$

or $f(x_i) - p_N^*(f)(x_i) = (-1)^{i+1} \|f - p_N^*(f)\|_\infty, \quad 0 \leq i \leq N + 1$

Corollary: $p_N^*(f)$ interpolates f in $N + 1$ points.

Illustration of Chebyshev's alternant theorem

$$N = 1$$

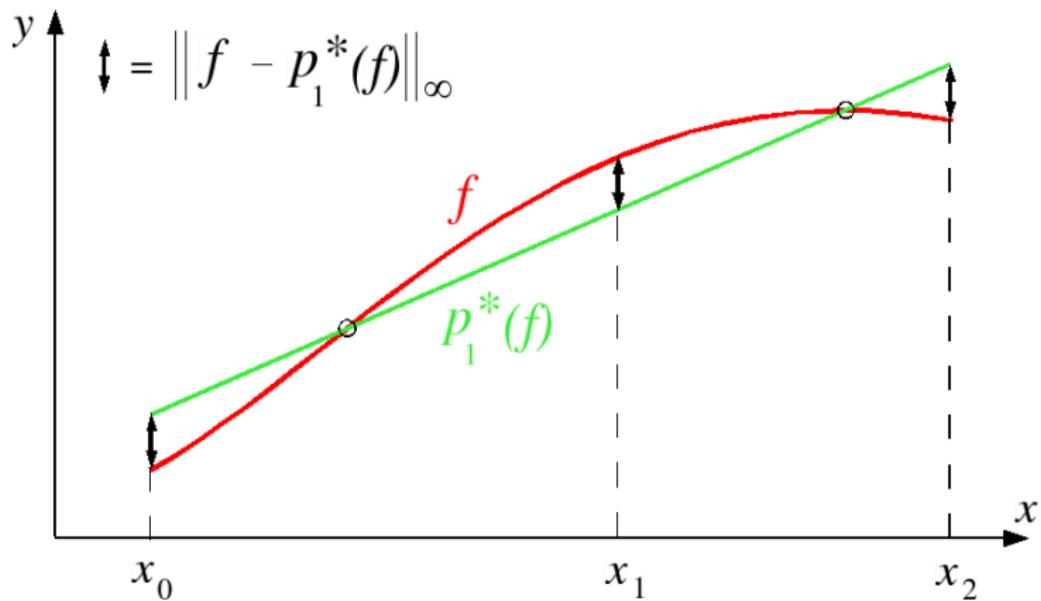
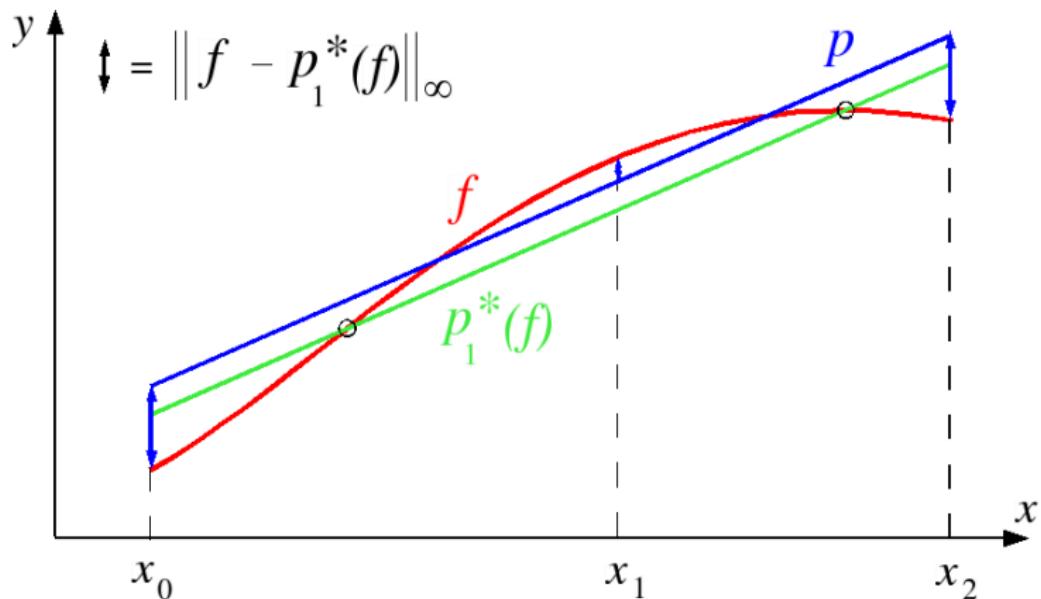


Illustration of Chebyshev's alternant theorem

$N = 1$



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Interpolation on an arbitrary grid

Definition: given an integer $N \geq 1$, a **grid** is a set of $N + 1$ points

$X = (x_i)_{0 \leq i \leq N}$ in $[-1,1]$ such that $-1 \leq x_0 < x_1 < \dots < x_N \leq 1$. The $N + 1$ points $(x_i)_{0 \leq i \leq N}$ are called the **nodes** of the grid.

Theorem

Given a function $f \in C^0([-1,1])$ and a grid of $N + 1$ nodes, $X = (x_i)_{0 \leq i \leq N}$, there exist a unique polynomial of degree N , $I_N^X f$, such that

$$I_N^X f(x_i) = f(x_i), \quad 0 \leq i \leq N$$

$I_N^X f$ is called the **interpolant** (or the **interpolating polynomial**) of f through the grid X .

Lagrange form of the interpolant

The interpolant $I_N^X f$ can be expressed in the *Lagrange form*:

$$I_N^X f(x) = \sum_{i=0}^N f(x_i) \ell_i^X(x),$$

where $\ell_i^X(x)$ is the i -th **Lagrange cardinal polynomial** associated with the grid X :

$$\ell_i^X(x) := \prod_{\substack{j=0 \\ j \neq i}}^N \frac{x - x_j}{x_i - x_j}, \quad 0 \leq i \leq N$$

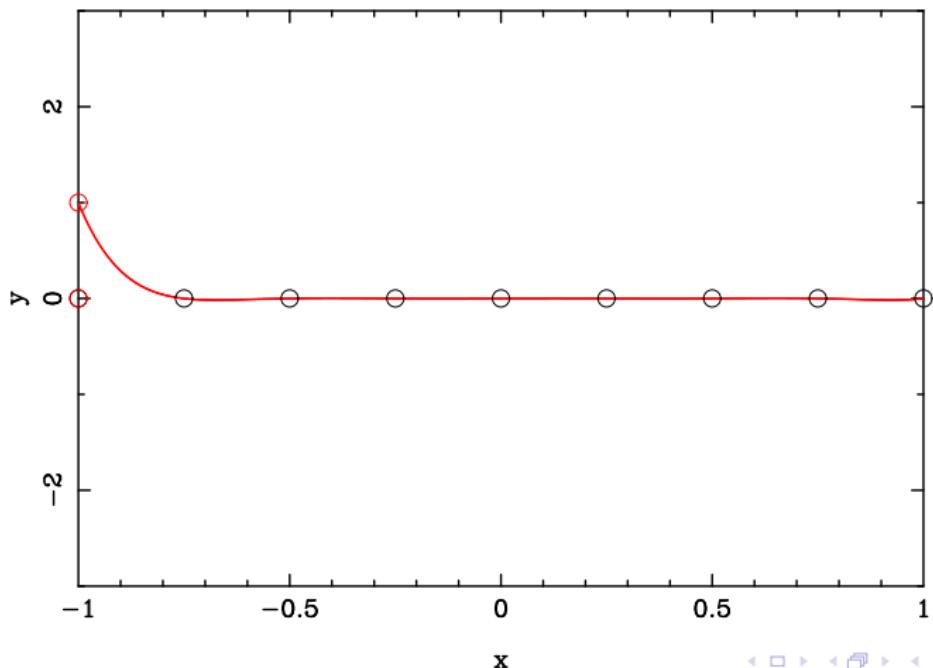
The Lagrange cardinal polynomials are such that

$$\ell_i^X(x_j) = \delta_{ij}, \quad 0 \leq i, j \leq N$$

Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_0^X(x)$

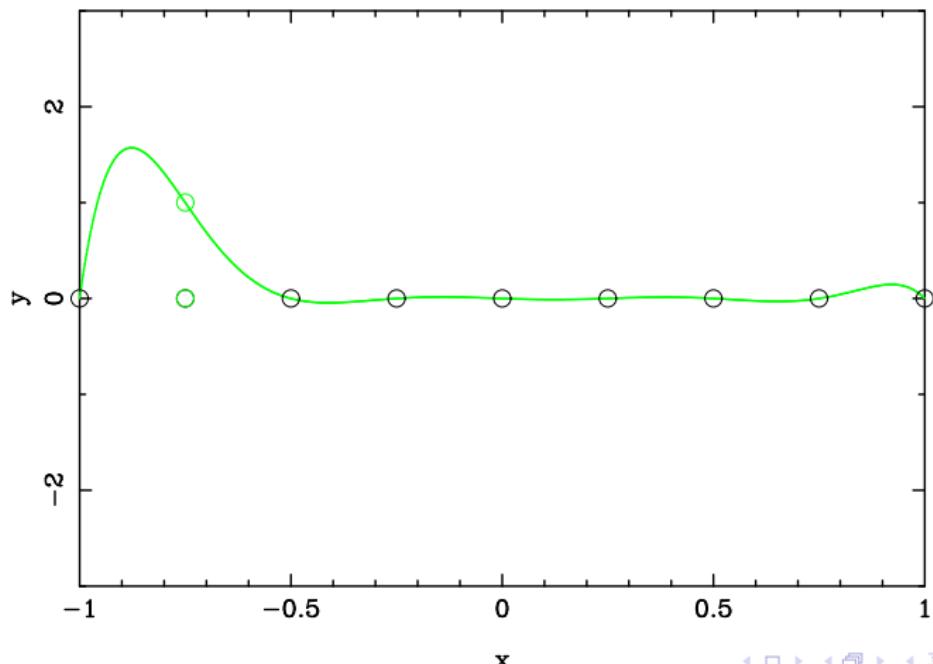
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_1^X(x)$

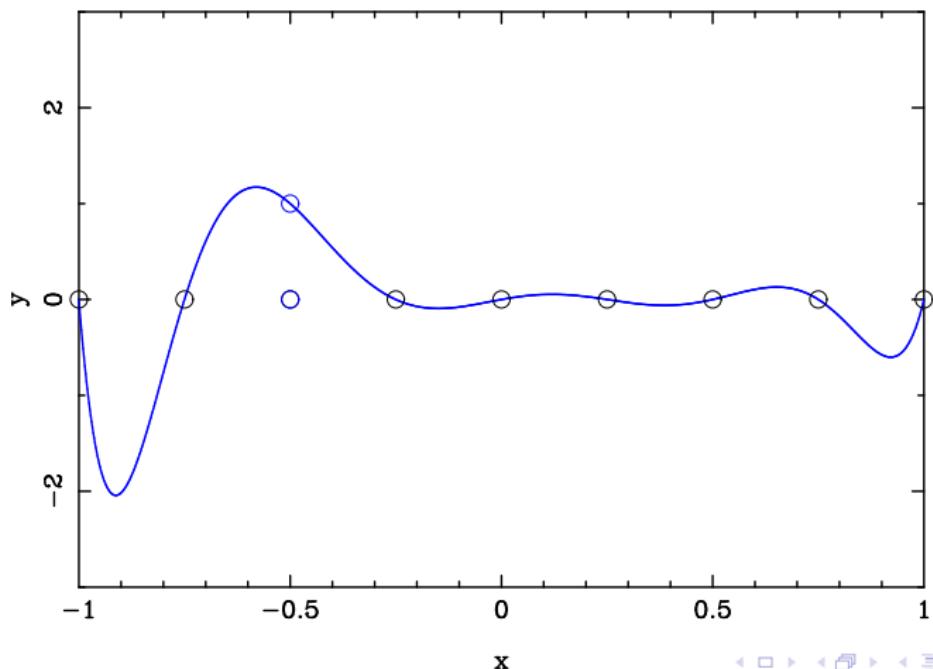
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_2^X(x)$

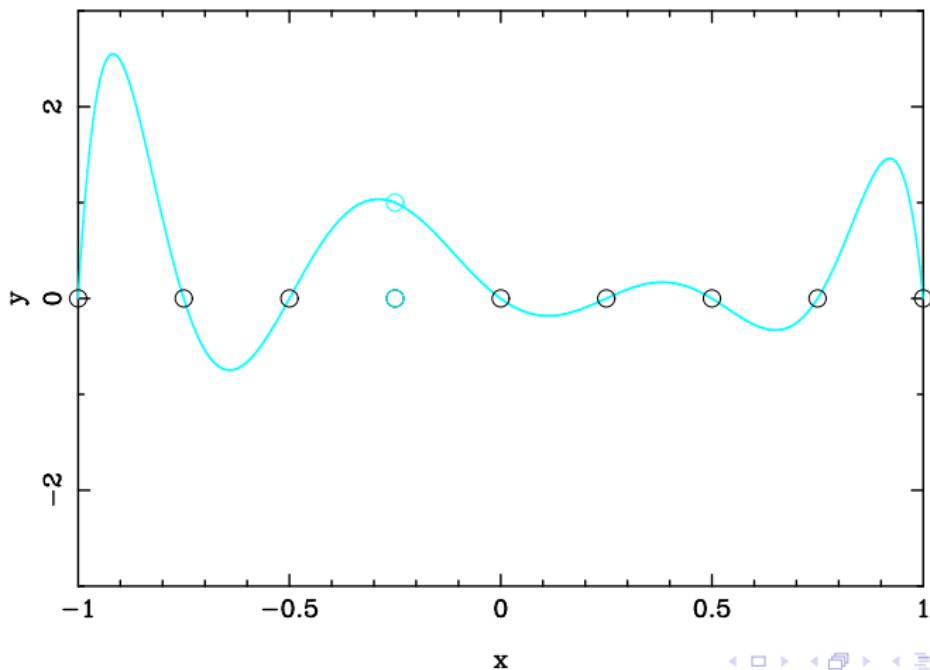
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_3^X(x)$

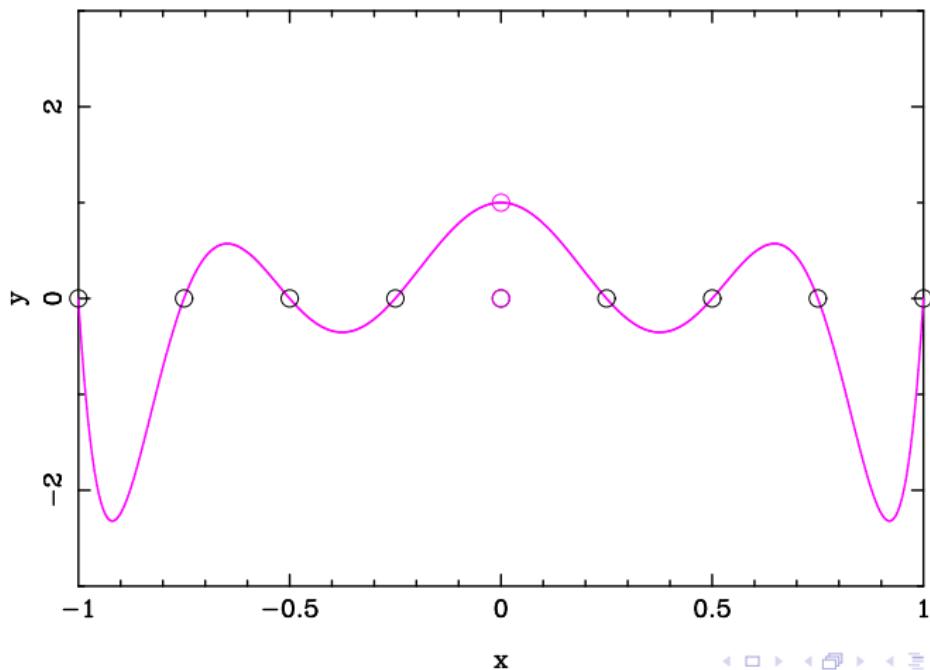
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_4^X(x)$

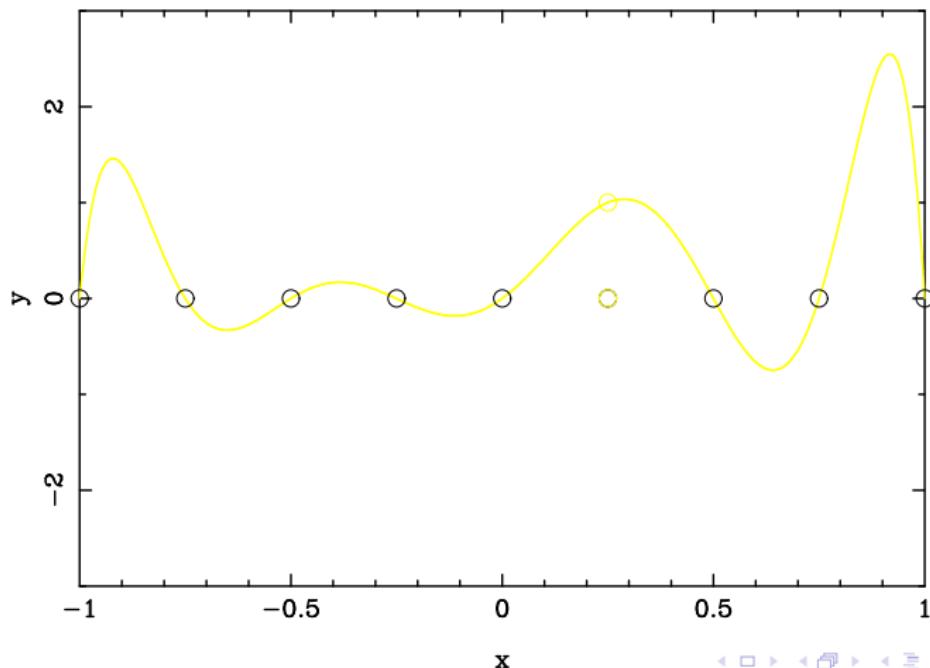
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_5^X(x)$

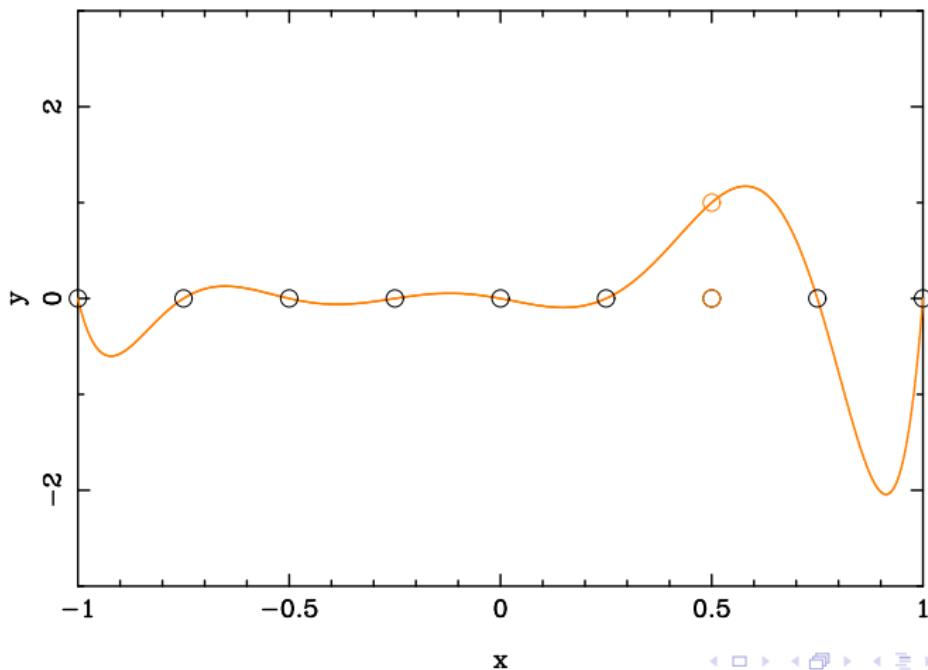
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_6^X(x)$

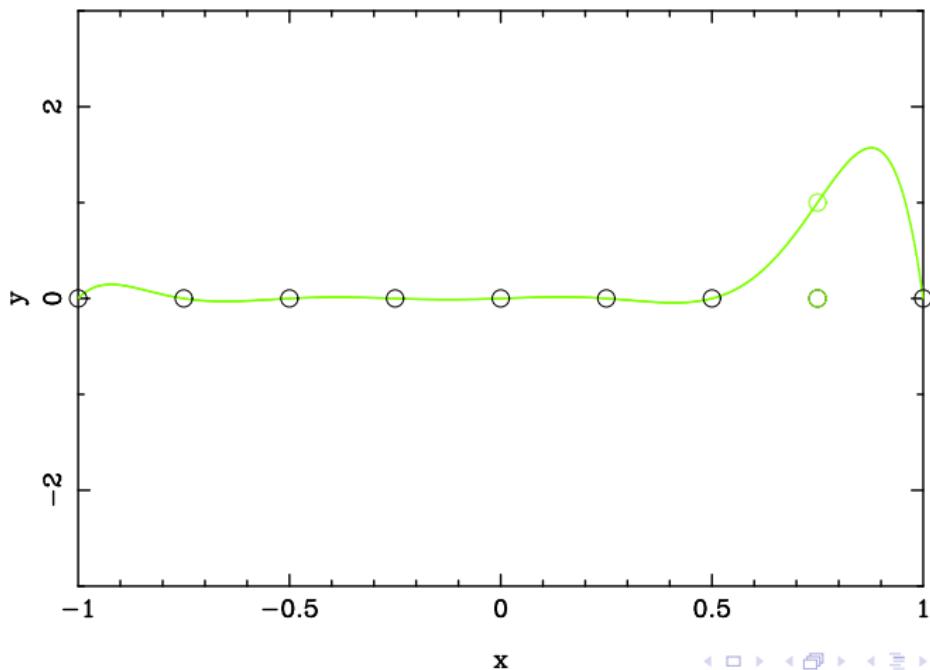
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_7^X(x)$

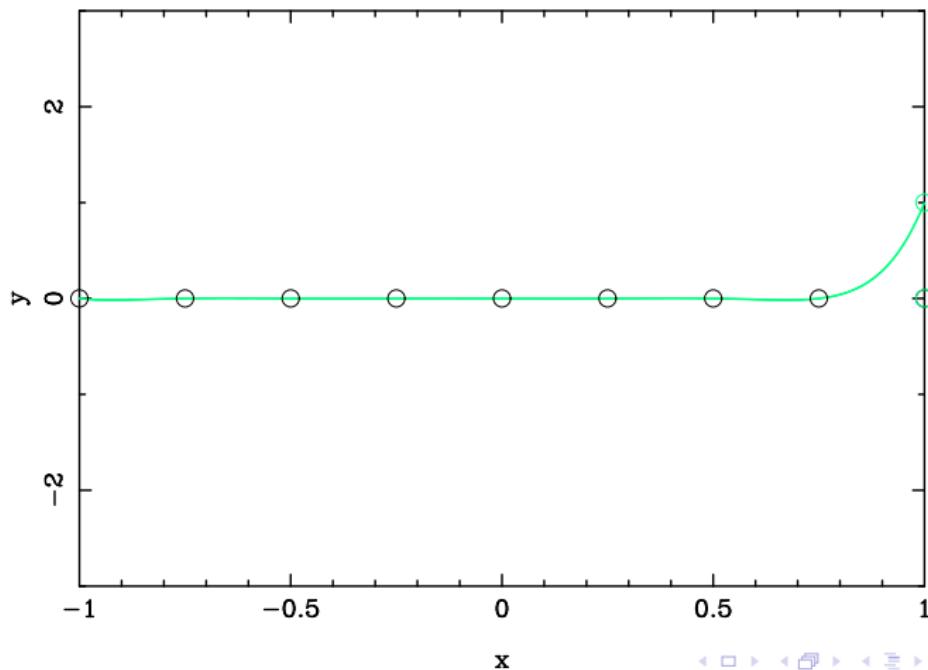
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$ $\ell_8^X(x)$

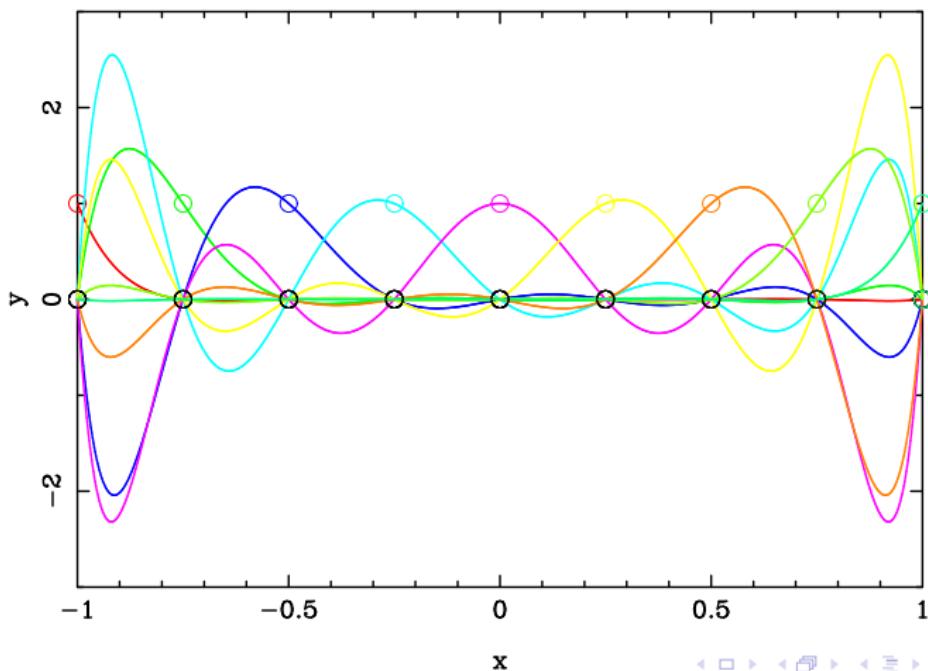
Lagrange polynomials



Examples of Lagrange polynomials

Uniform grid $N = 8$

Lagrange polynomials



Interpolation error with respect to the best approximation error

Let $N \in \mathbb{N}$, $X = (x_i)_{0 \leq i \leq N}$ a grid of $N + 1$ nodes and $f \in C^0([-1, 1])$.

Let us consider the interpolant $I_N^X f$ of f through the grid X .

The best approximation polynomial $p_N^*(f)$ is also an interpolant of f at $N + 1$ nodes (in general different from X)

◀ reminder

How does the error $\|f - I_N^X f\|_\infty$ behave with respect to the smallest possible error $\|f - p_N^*(f)\|_\infty$?

The answer is given by the formula:

$$\|f - I_N^X f\|_\infty \leq (1 + \Lambda_N(X)) \|f - p_N^*(f)\|_\infty$$

where $\Lambda_N(X)$ is the **Lebesgue constant** relative to the grid X :

$$\Lambda_N(X) := \max_{x \in [-1, 1]} \sum_{i=0}^N |\ell_i^X(x)|$$

Lebesgue constant

The Lebesgue constant contains all the information on the effects of the choice of X on $\|f - I_N^X f\|_\infty$.

Theorem (Erdős, 1961)

For any choice of the grid X , there exists a constant $C > 0$ such that

$$\Lambda_N(X) > \frac{2}{\pi} \ln(N+1) - C$$

Corollary: $\Lambda_N(X) \rightarrow \infty$ as $N \rightarrow \infty$

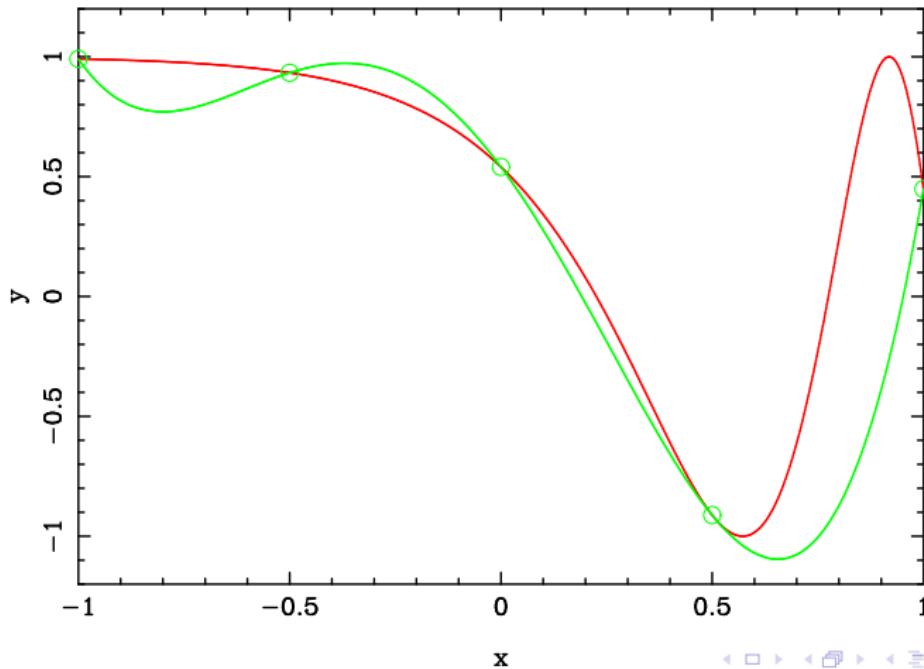
In particular, for a **uniform** grid, $\Lambda_N(X) \sim \frac{2^{N+1}}{eN \ln N}$ as $N \rightarrow \infty$!

This means that for any choice of type of sampling of $[-1, 1]$, there exists a continuous function $f \in C^0([-1, 1])$ such that $I_N^X f$ does not converge uniformly towards f .

Example: uniform interpolation of a “gentle” function

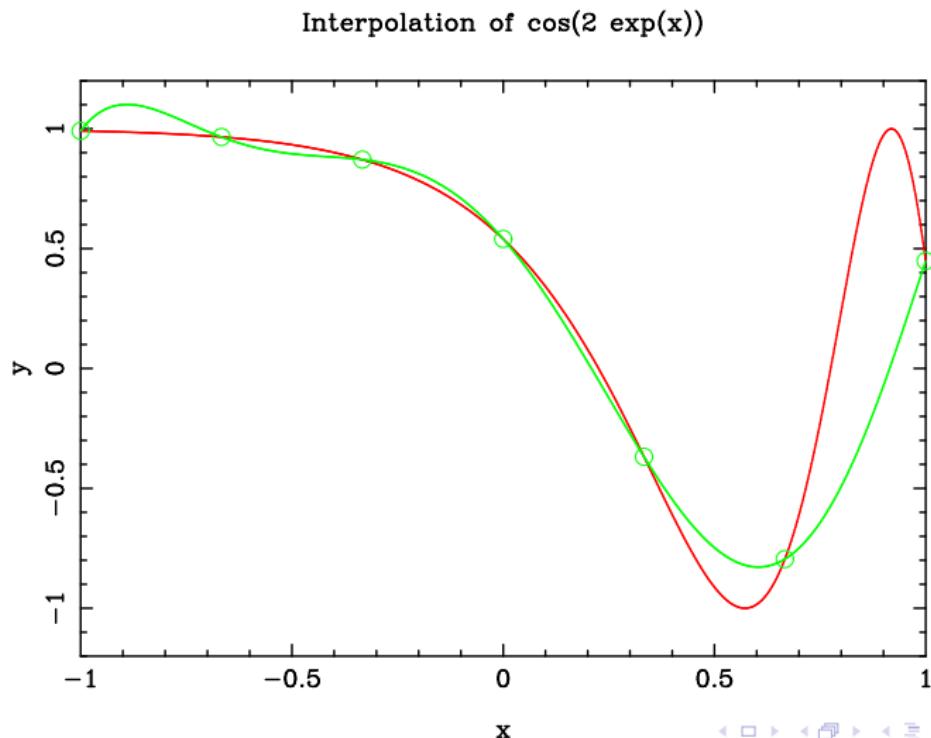
$f(x) = \cos(2 \exp(x))$ uniform grid $N = 4$: $\|f - I_4^X f\|_\infty \simeq 1.40$

Interpolation of $\cos(2 \exp(x))$



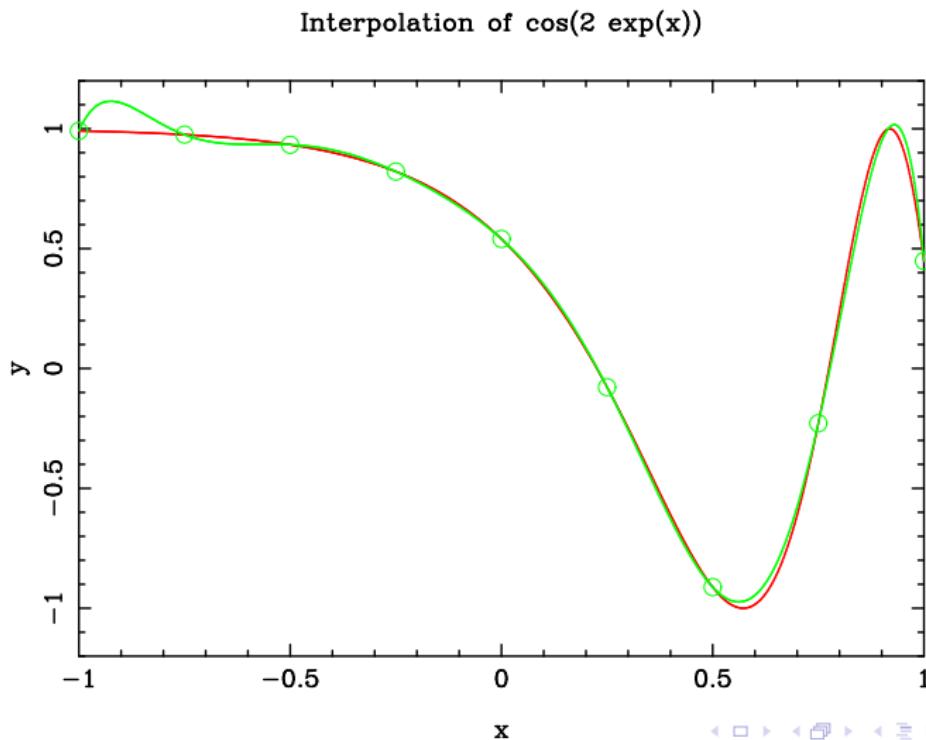
Example: uniform interpolation of a “gentle” function

$f(x) = \cos(2 \exp(x))$ uniform grid $N = 6$: $\|f - I_6^X f\|_\infty \simeq 1.05$



Example: uniform interpolation of a “gentle” function

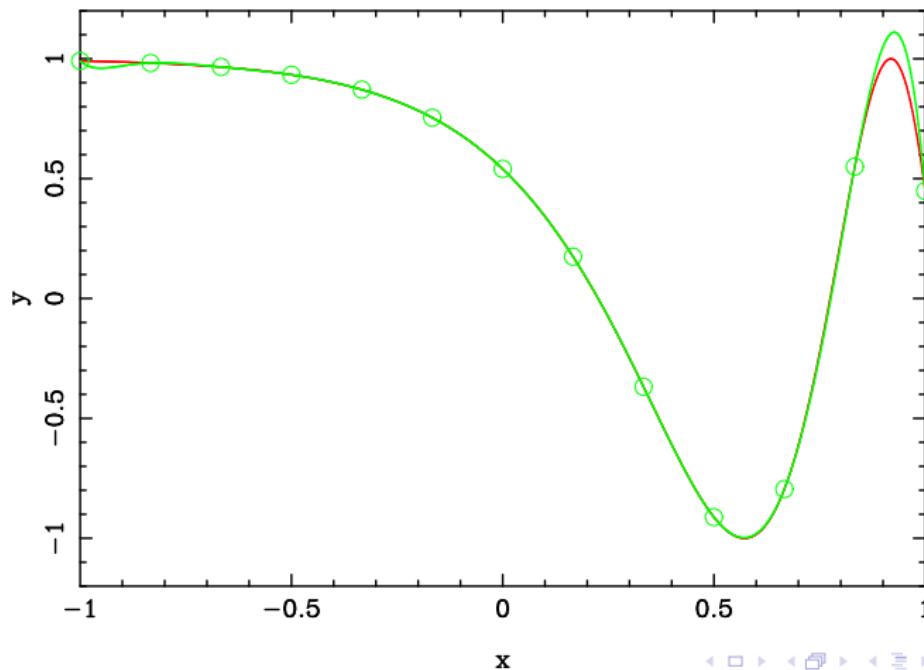
$f(x) = \cos(2 \exp(x))$ uniform grid $N = 8$: $\|f - I_8^X f\|_\infty \simeq 0.13$



Example: uniform interpolation of a “gentle” function

$f(x) = \cos(2 \exp(x))$ uniform grid $N = 12$: $\|f - I_{12}^X f\|_\infty \simeq 0.13$

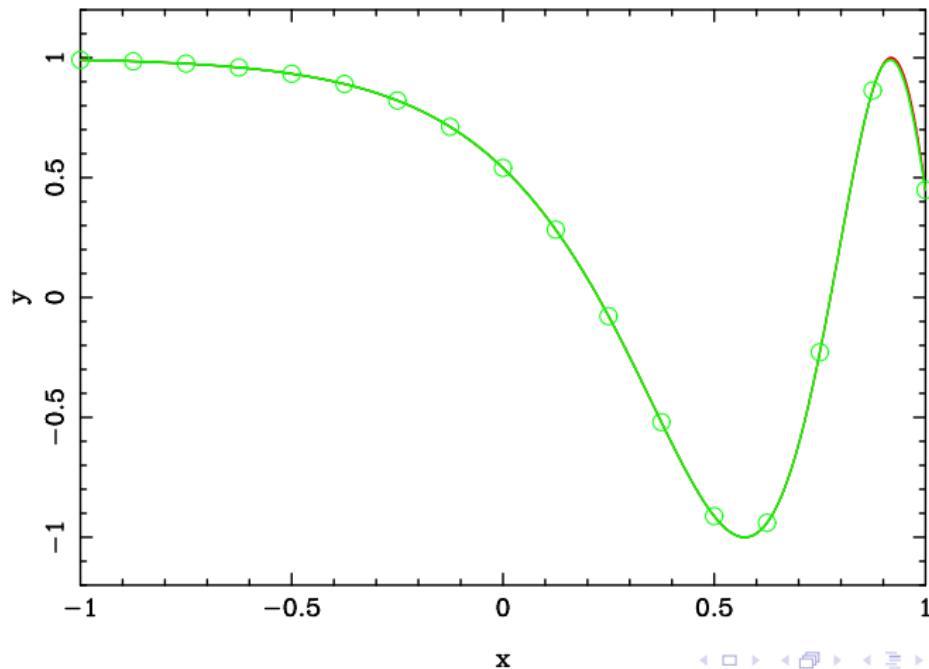
Interpolation of $\cos(2 \exp(x))$



Example: uniform interpolation of a “gentle” function

$f(x) = \cos(2 \exp(x))$ uniform grid $N = 16$: $\|f - I_{16}^X f\|_\infty \simeq 0.025$

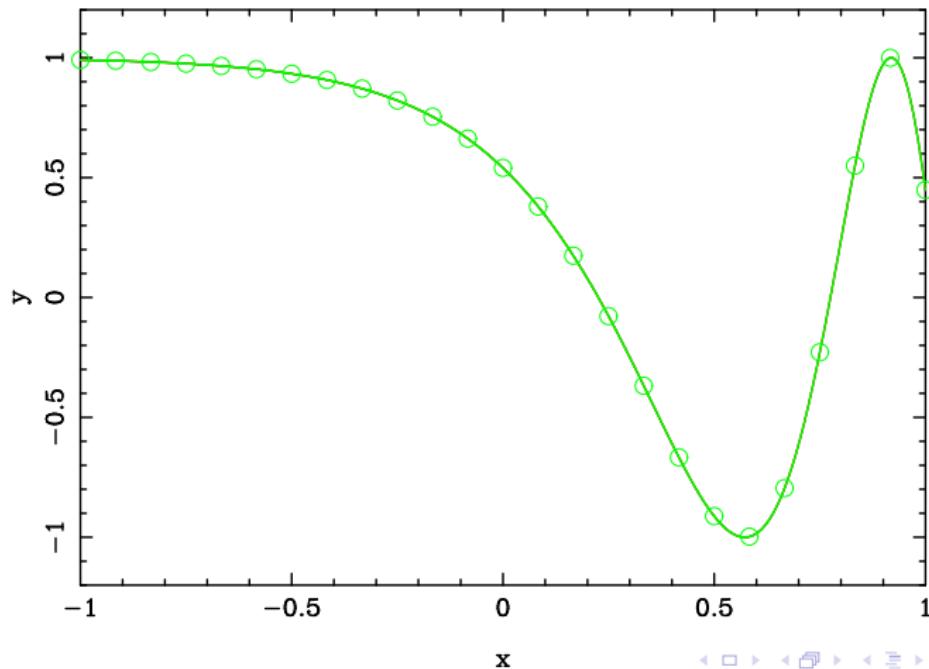
Interpolation of $\cos(2 \exp(x))$



Example: uniform interpolation of a “gentle” function

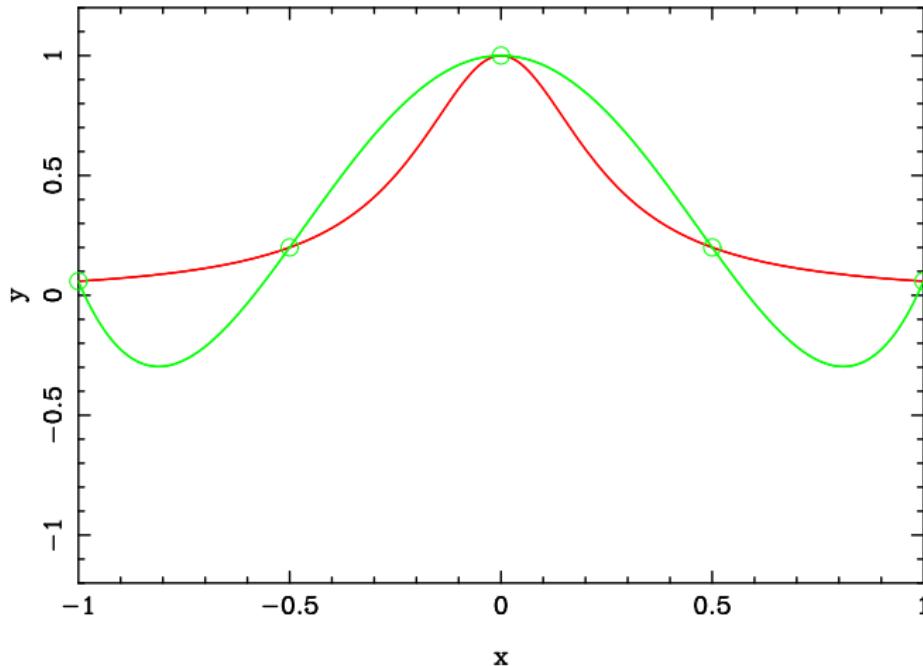
$f(x) = \cos(2 \exp(x))$ uniform grid $N = 24$: $\|f - I_{24}^X f\|_\infty \simeq 4.6 \cdot 10^{-4}$

Interpolation of $\cos(2 \exp(x))$



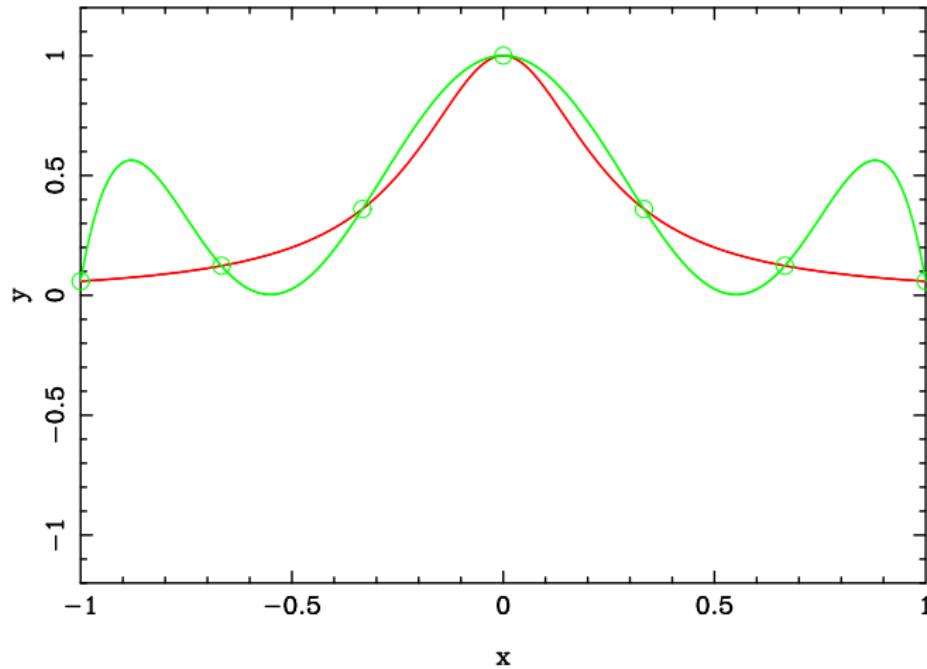
Runge phenomenon

$$f(x) = \frac{1}{1 + 16x^2} \quad \text{uniform grid } N = 4 : \|f - I_4^X f\|_\infty \simeq 0.39$$



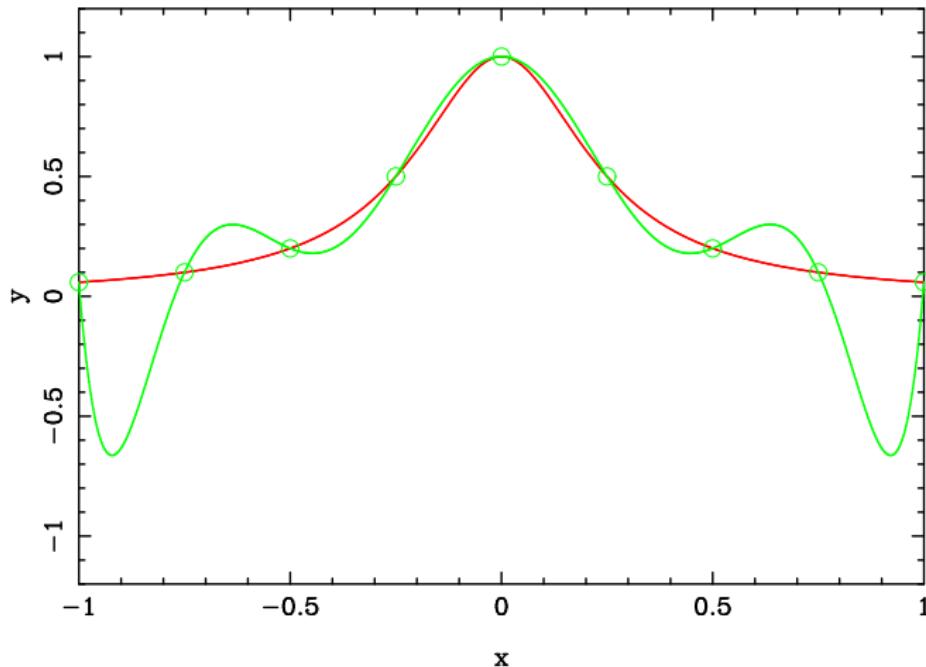
Runge phenomenon

$$f(x) = \frac{1}{1 + 16x^2} \quad \text{uniform grid } N = 6 : \|f - I_6^X f\|_\infty \simeq 0.49$$



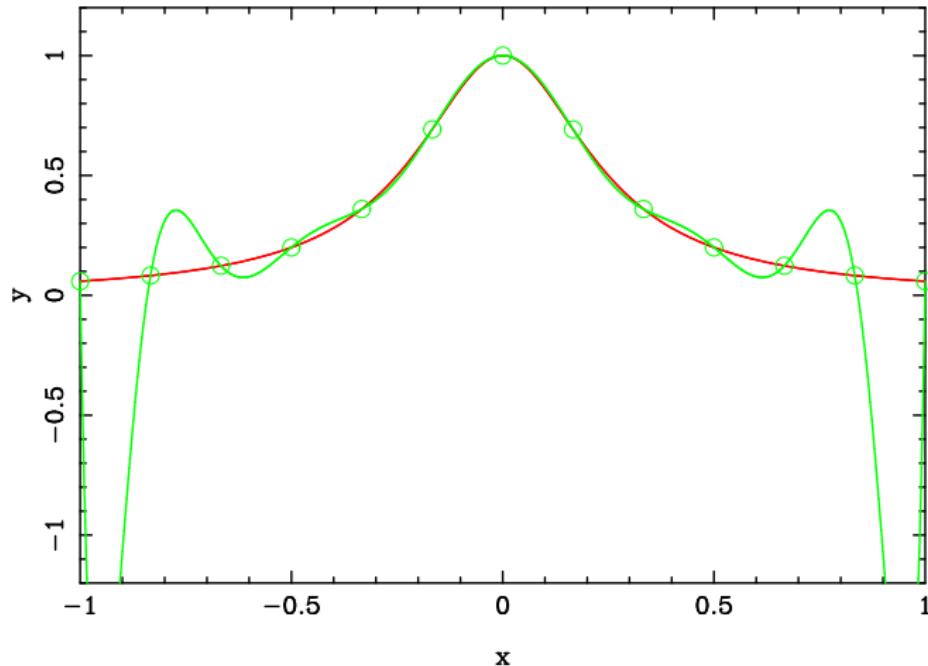
Runge phenomenon

$$f(x) = \frac{1}{1 + 16x^2} \quad \text{uniform grid } N = 8 : \|f - I_8^X f\|_\infty \simeq 0.73$$



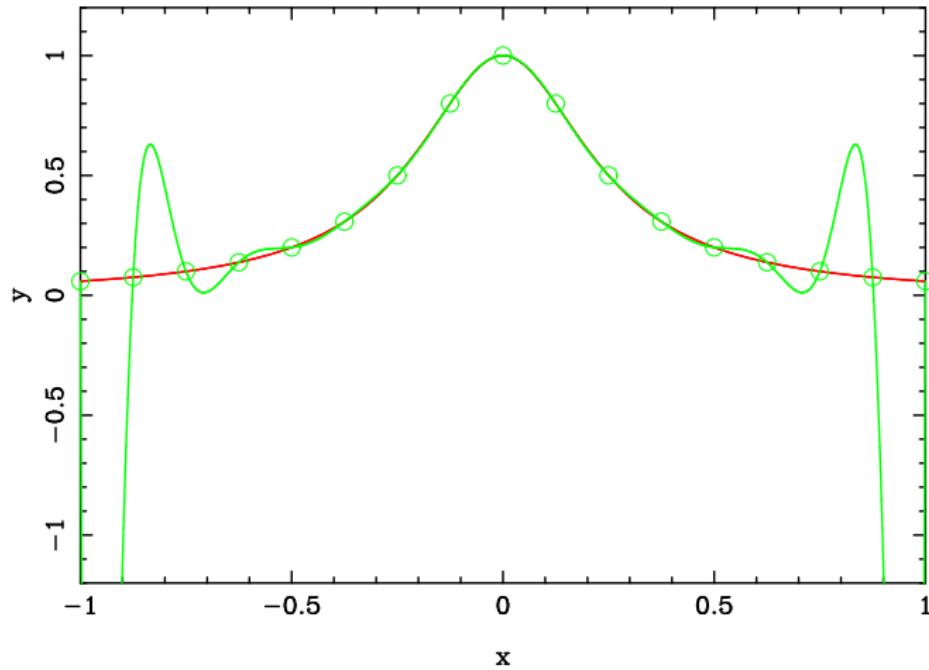
Runge phenomenon

$$f(x) = \frac{1}{1+16x^2} \quad \text{uniform grid } N = 12 : \|f - I_{12}^X f\|_\infty \simeq 1.97$$



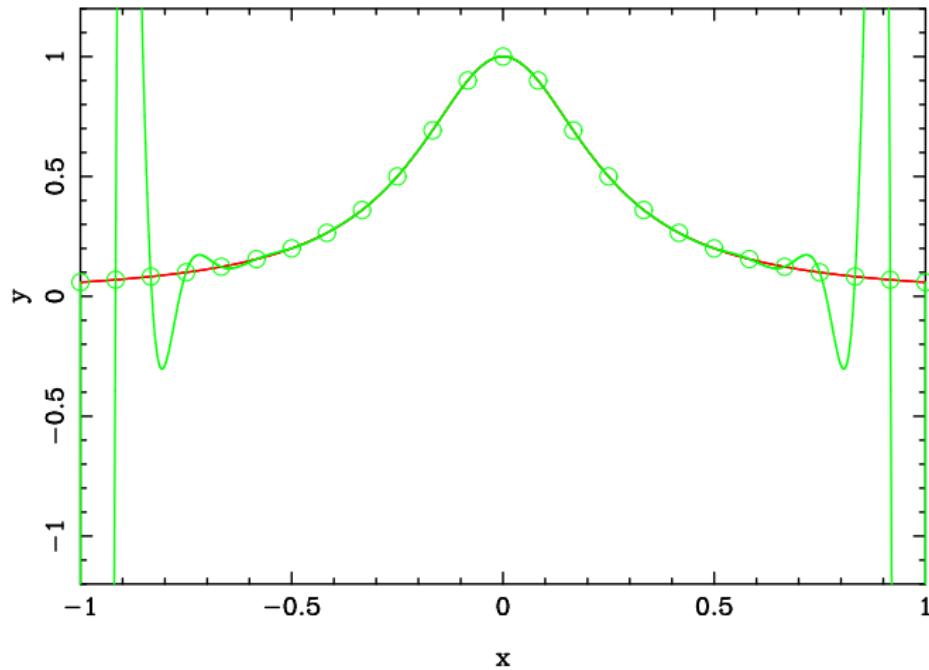
Runge phenomenon

$$f(x) = \frac{1}{1 + 16x^2} \quad \text{uniform grid } N = 16 : \|f - I_{16}^X f\|_\infty \simeq 5.9$$



Runge phenomenon

$$f(x) = \frac{1}{1 + 16x^2} \quad \text{uniform grid } N = 24 : \|f - I_{24}^X f\|_\infty \simeq 62$$



Evaluation of the interpolation error

Let us assume that the function f is sufficiently smooth to have derivatives at least up to the order $N + 1$, with $f^{(N+1)}$ continuous, i.e. $f \in C^{N+1}([-1, 1])$.

Theorem (Cauchy)

If $f \in C^{N+1}([-1, 1])$, then for any grid X of $N + 1$ nodes, and for any $x \in [-1, 1]$, the interpolation error at x is

$$f(x) - I_N^X(x) = \frac{f^{(N+1)}(\xi)}{(N+1)!} \omega_{N+1}^X(x) \quad (1)$$

where $\xi = \xi(x) \in [-1, 1]$ and $\omega_{N+1}^X(x)$ is the nodal polynomial associated with the grid X .

Definition: The **nodal polynomial** associated with the grid X is the unique polynomial of degree $N + 1$ and leading coefficient 1 whose zeros are the $N + 1$ nodes of X :

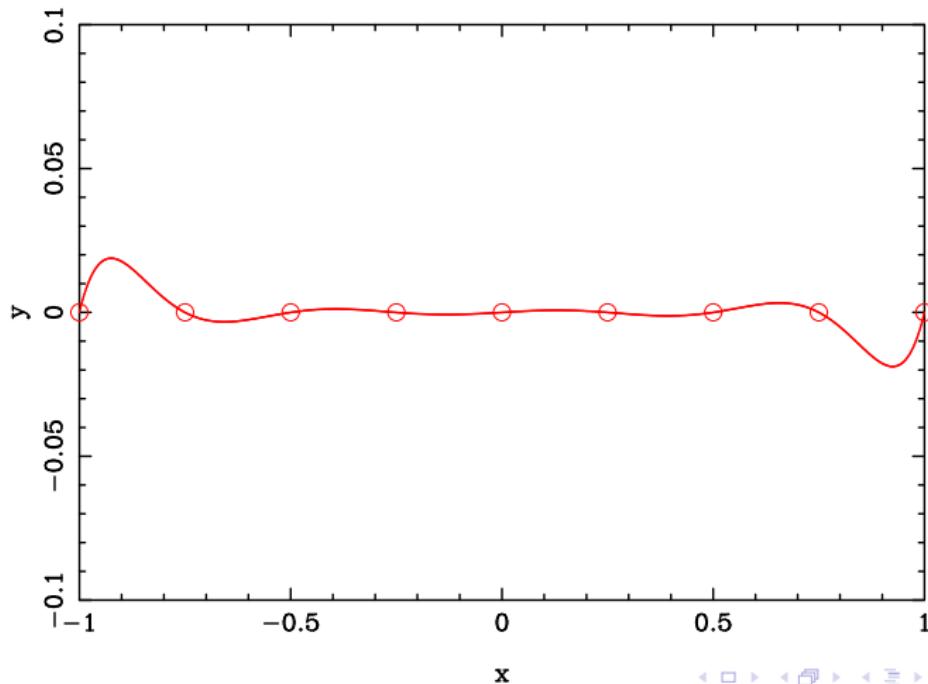
$$\omega_{N+1}^X(x) := \prod_{i=0}^N (x - x_i)$$

Example of nodal polynomial

Uniform grid

$N = 8$

Nodal polynomial



Minimizing the interpolation error by the choice of grid

In Eq. (1), we have no control on $f^{(N+1)}$, which can be large.

For example, for $f(x) = 1/(1 + \alpha^2 x^2)$, $\|f^{(N+1)}\|_\infty = (N+1)! \alpha^{N+1}$.

Idea: choose the grid X so that $\omega_{N+1}^X(x)$ is small, i.e. $\|\omega_{N+1}^X\|_\infty$ is small.

Notice: $\omega_{N+1}^X(x)$ has leading coefficient 1: $\omega_{N+1}^X(x) = x^{N+1} + \sum_{i=0}^N a_i x^i$.

Theorem (Chebyshev)

Among all the polynomials of degree $N+1$ and leading coefficient 1, the unique polynomial which has the smallest uniform norm on $[-1, 1]$ is the $(N+1)$ -th Chebyshev polynomial divided by 2^N : $T_{N+1}(x)/2^N$.

Since $\|T_{N+1}\|_\infty = 1$, we conclude that if we choose the grid nodes $(x_i)_{0 \leq i \leq N}$ to be the $N+1$ zeros of the Chebyshev polynomial T_{N+1} , we have

$$\|\omega_{N+1}^X\|_\infty = \frac{1}{2^N}$$

and this is the smallest possible value.

Chebyshev-Gauss grid

The grid $X = (x_i)_{0 \leq i \leq N}$ such that the x_i 's are the $N + 1$ zeros of the Chebyshev polynomial of degree $N + 1$ is called the **Chebyshev-Gauss (CG) grid**.

It has much better interpolation properties than the uniform grid considered so far. In particular, from Eq. (1), for any function $f \in C^{N+1}([-1, 1])$,

$$\|f - I_N^{\text{CG}} f\|_\infty \leq \frac{1}{2^N (N+1)!} \|f^{(N+1)}\|_\infty$$

If $f^{(N+1)}$ is uniformly bounded, the convergence of the interpolant $I_N^{\text{CG}} f$ towards f when $N \rightarrow \infty$ is then extremely fast.

Also the Lebesgue constant associated with the Chebyshev-Gauss grid is small:

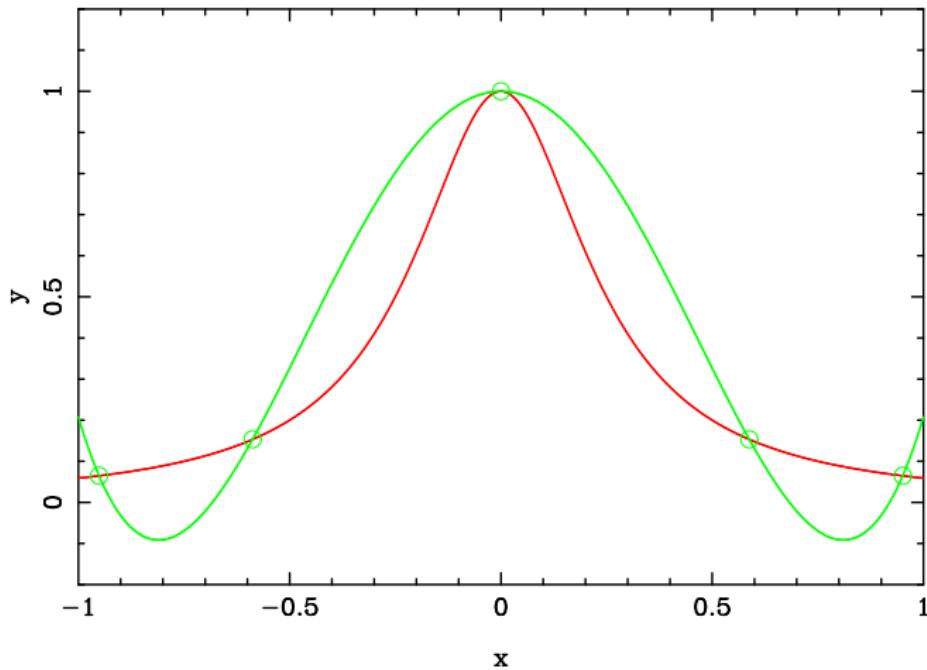
$$\Lambda_N(\text{CG}) \sim \frac{2}{\pi} \ln(N+1) \quad \text{as } N \rightarrow \infty$$

This is much better than uniform grids and close to the optimal value

reminder

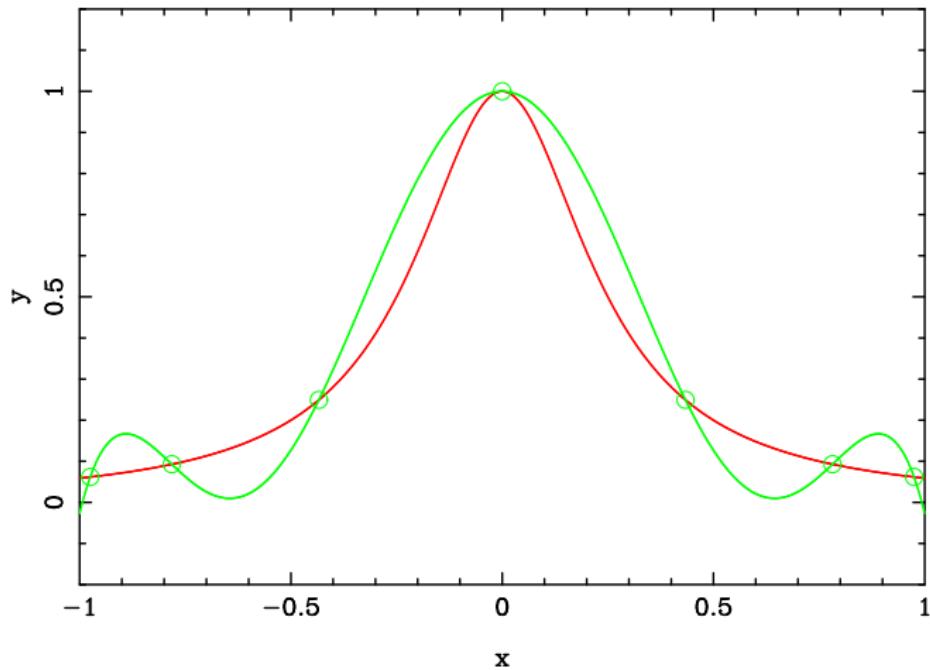
Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

$$f(x) = \frac{1}{1+16x^2} \quad \text{CG grid } N=4 : \|f - I_4^{\text{CG}} f\|_\infty \simeq 0.31$$



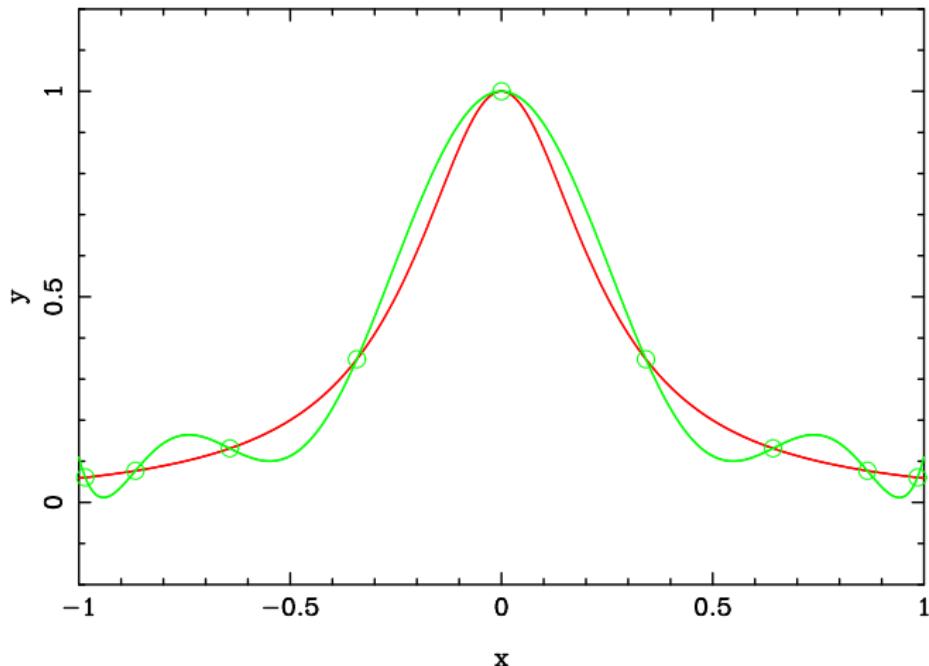
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$$f(x) = \frac{1}{1+16x^2} \quad \text{CG grid } N=6 : \|f - I_6^{\text{CG}} f\|_\infty \simeq 0.18$$



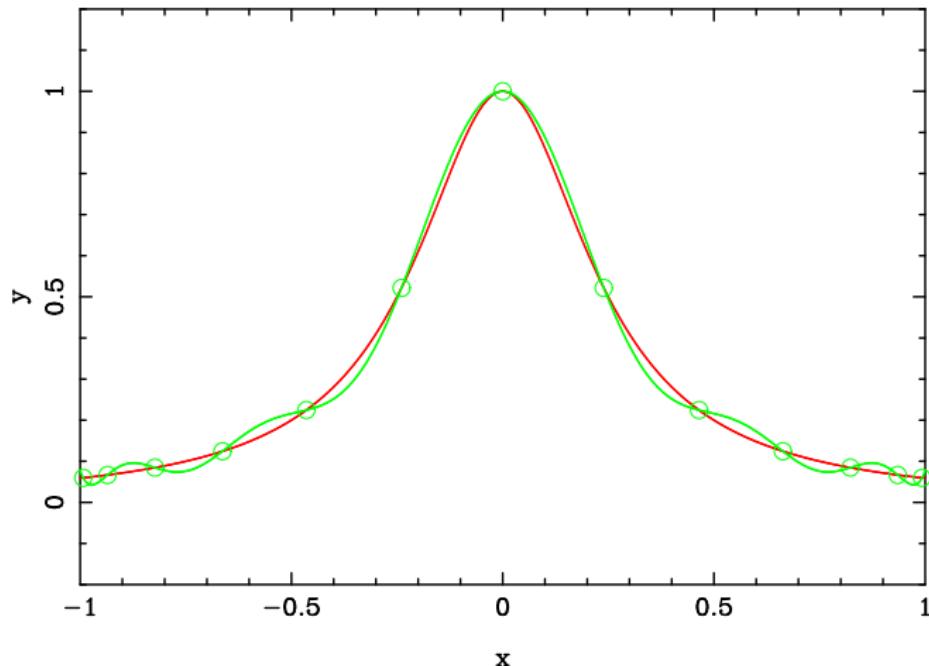
Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

$$f(x) = \frac{1}{1+16x^2} \quad \text{CG grid } N=8 : \|f - I_8^{\text{CG}} f\|_\infty \simeq 0.10$$



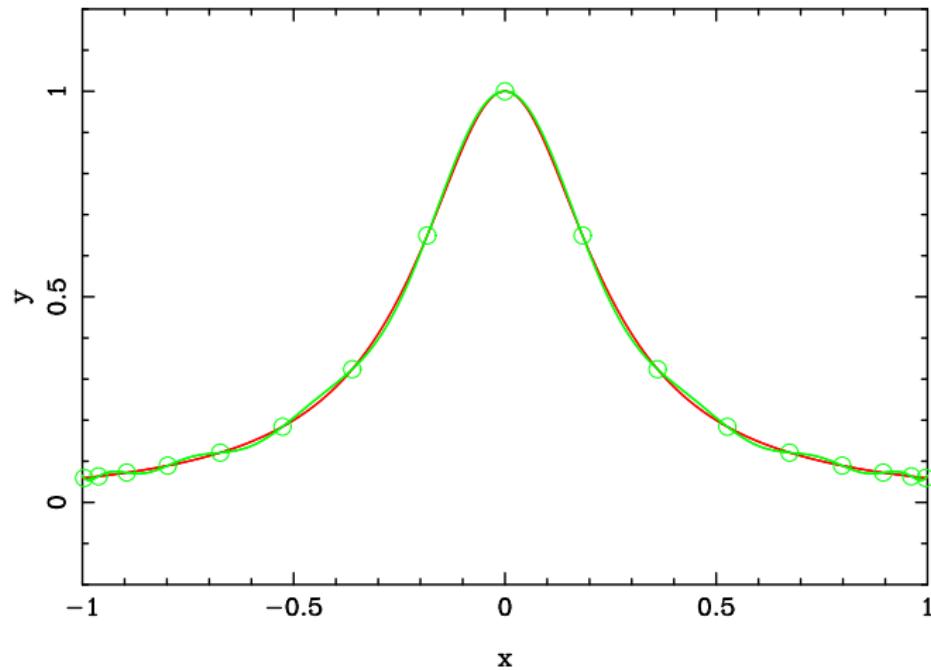
Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

$$f(x) = \frac{1}{1+16x^2} \quad \text{CG grid } N = 12 : \|f - I_{12}^{\text{CG}} f\|_\infty \simeq 3.8 \cdot 10^{-2}$$



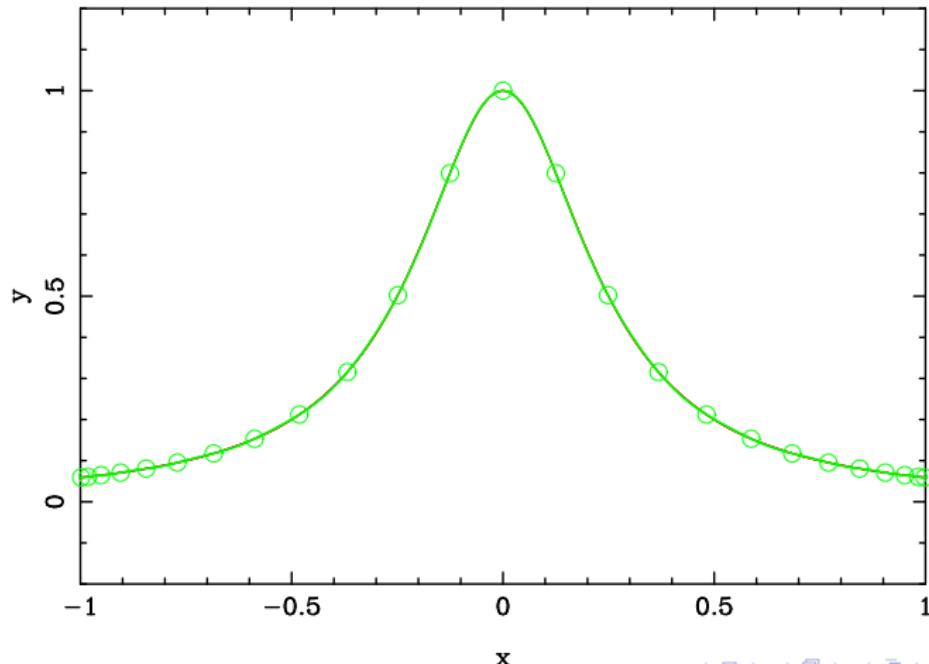
Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

$$f(x) = \frac{1}{1+16x^2} \quad \text{CG grid } N = 16 : \|f - I_{16}^{\text{CG}} f\|_\infty \simeq 1.5 \cdot 10^{-2}$$



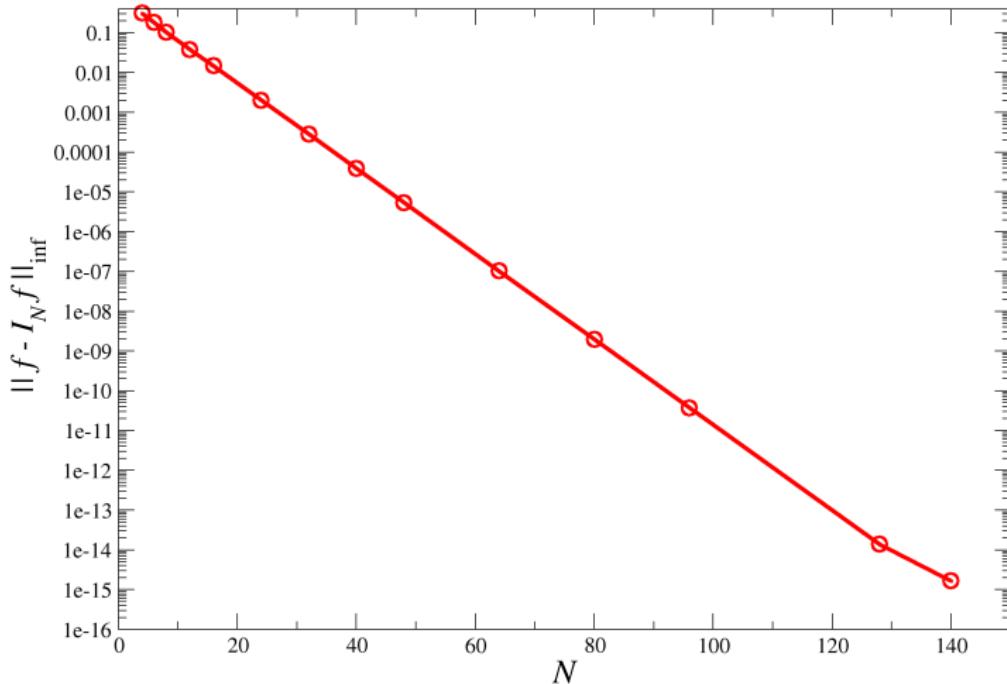
Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

$f(x) = \frac{1}{1+16x^2}$ CG grid $N = 24$: $\|f - I_{24}^{\text{CG}} f\|_\infty \simeq 2.0 \cdot 10^{-3}$
no Runge phenomenon !



Example: Chebyshev-Gauss interpolation of $f(x) = \frac{1}{1+16x^2}$

Variation of the interpolation error as N increases



Chebyshev polynomials = orthogonal polynomials

The Chebyshev polynomials, the zeros of which provide the Chebyshev-Gauss nodes, constitute a family of **orthogonal polynomials**, and the Chebyshev-Gauss nodes are associated to **Gauss quadratures**.

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Hilbert space $L_w^2(-1, 1)$

Framework: Let us consider the functional space

$$L_w^2(-1, 1) = \left\{ f : (-1, 1) \rightarrow \mathbb{R}, \int_{-1}^1 f(x)^2 w(x) dx < \infty \right\}$$

where $w : (-1, 1) \rightarrow (0, \infty)$ is an integrable function, called the **weight function**.

$L_w^2(-1, 1)$ is a **Hilbert space** for the scalar product

$$(f|g)_w := \int_{-1}^1 f(x) g(x) w(x) dx$$

with the associated norm

$$\|f\|_w := (f|f)_w^{1/2}$$

Orthogonal polynomials

The set \mathbb{P} of polynomials on $[-1, 1]$ is a subspace of $L_w^2(-1, 1)$.
 A family of **orthogonal polynomials** is a set $(p_i)_{i \in \mathbb{N}}$ such that

- $p_i \in \mathbb{P}$
- $\deg p_i = i$
- $i \neq j \Rightarrow (p_i | p_j)_w = 0$

$(p_i)_{i \in \mathbb{N}}$ is then a basis of the vector space \mathbb{P} : $\mathbb{P} = \text{span} \{p_i, i \in \mathbb{N}\}$

Theorem

A family of orthogonal polynomial $(p_i)_{i \in \mathbb{N}}$ is a **Hilbert basis** of $L_w^2(-1, 1)$:

$$\forall f \in L_w^2(-1, 1), \quad f = \sum_{i=0}^{\infty} \tilde{f}_i p_i \quad \text{with } \tilde{f}_i := \frac{(f | p_i)_w}{\|p_i\|_w^2}.$$

The above infinite sum means $\lim_{N \rightarrow \infty} \left\| f - \sum_{i=0}^N \tilde{f}_i p_i \right\|_w = 0$

Jacobi polynomials

Jacobi polynomials are orthogonal polynomials with respect to the weight

$$w(x) = (1-x)^\alpha (1+x)^\beta$$

Subcases:

- Legendre polynomials $P_n(x)$: $\alpha = \beta = 0$, i.e. $w(x) = 1$
- Chebyshev polynomials $T_n(x)$: $\alpha = \beta = -\frac{1}{2}$, i.e. $w(x) = \frac{1}{\sqrt{1-x^2}}$

Jacobi polynomials are eigenfunctions of the singular¹ **Sturm-Liouville problem**

$$-\frac{d}{dx} \left[(1-x^2) w(x) \frac{du}{dx} \right] = \lambda w(x) u, \quad x \in (-1, 1)$$

¹singular means that the coefficient in front of du/dx vanishes at the extremities of the interval $[-1, 1]$

Legendre polynomials

$$w(x) = 1: \int_{-1}^1 P_i(x)P_j(x) dx = \frac{2}{2i+1} \delta_{ij}$$

$$P_0(x) = 1$$

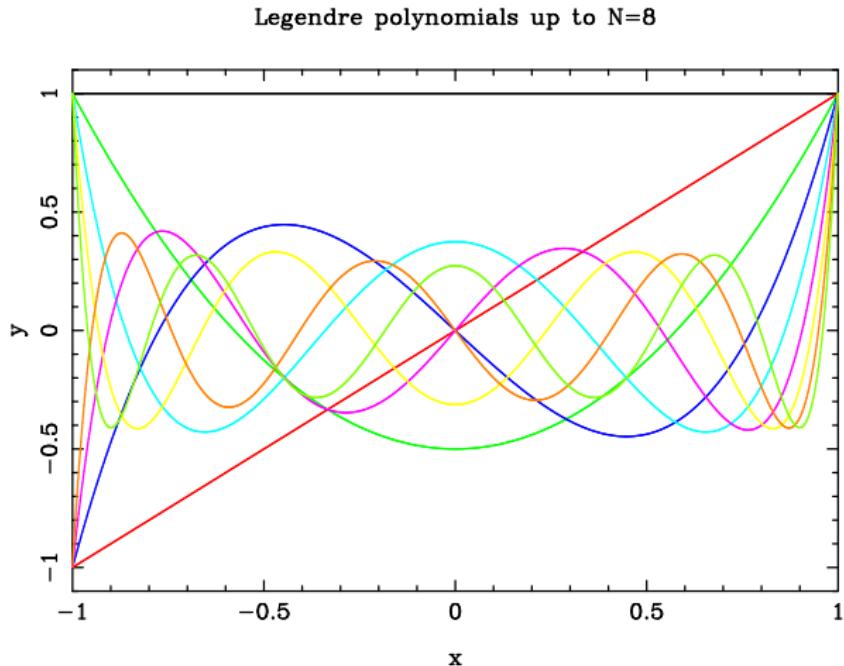
$$P_1(x) = x$$

$$P_2(x) = \frac{1}{2}(3x^2 - 1)$$

$$P_3(x) = \frac{1}{2}(5x^3 - 3x)$$

$$P_4(x) = \frac{1}{8}(35x^4 - 30x^2 + 3)$$

$$P_{i+1}(x) = \frac{2i+1}{i+1} x P_i(x) - \frac{i}{i+1} P_{i-1}(x)$$



Chebyshev polynomials

$$w(x) = \frac{1}{\sqrt{1-x^2}}: \quad \int_{-1}^1 T_i(x)T_j(x) \frac{dx}{\sqrt{1-x^2}} = \frac{\pi}{2}(1 + \delta_{0i}) \delta_{ij}$$

Chebyshev polynomials up to N=8

$$T_0(x) = 1$$

$$T_1(x) = x$$

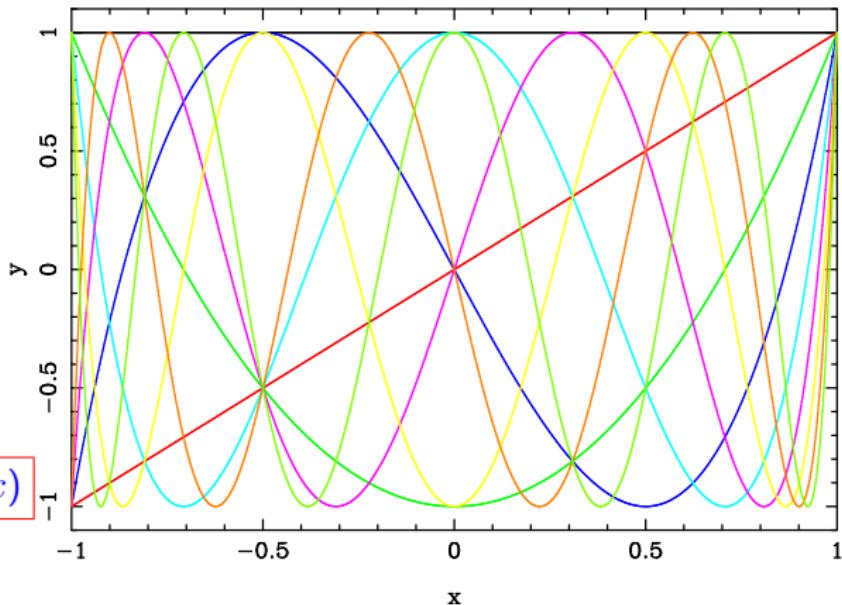
$$T_2(x) = 2x^2 - 1$$

$$T_3(x) = 4x^3 - 3x$$

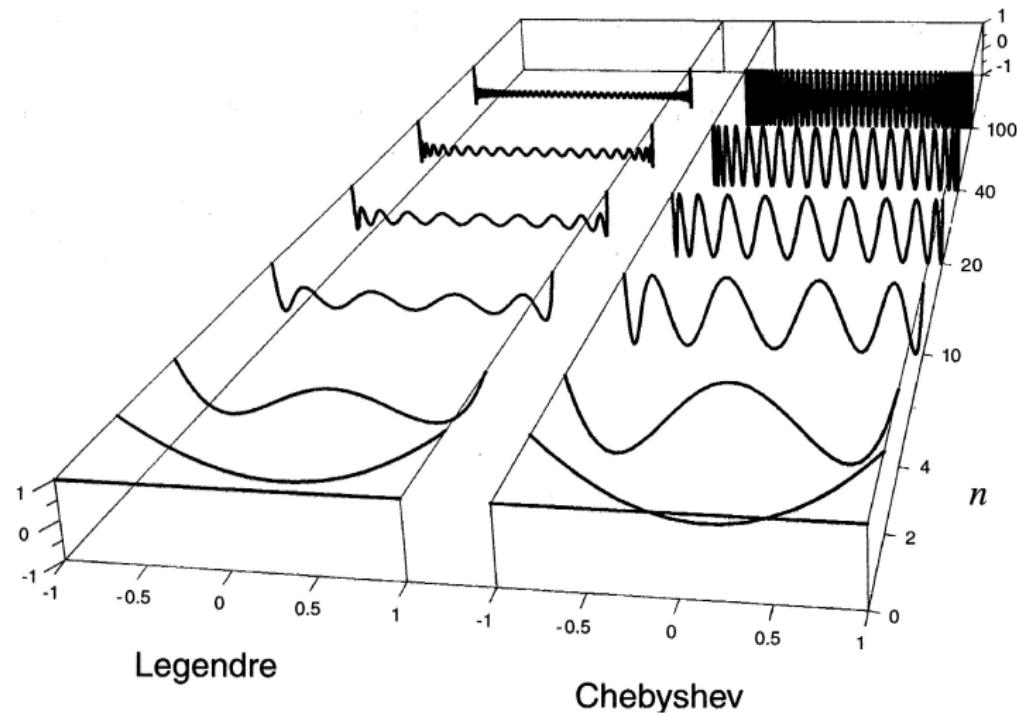
$$T_4(x) = 8x^4 - 8x^2 + 1$$

$$\cos(n\theta) = T_n(\cos \theta)$$

$$T_{i+1}(x) = 2xT_i(x) - T_{i-1}(x)$$



Legendre and Chebyshev compared



[from Fornberg (1998)]

Orthogonal projection on \mathbb{P}_N

Let us consider $f \in L_w^2(-1, 1)$ and a family $(p_i)_{i \in \mathbb{N}}$ of orthogonal polynomials with respect to the weight w .

Since $(p_i)_{i \in \mathbb{N}}$ is a Hilbert basis of $L_w^2(-1, 1)$ ◀ reminder

we have $f(x) = \sum_{i=0}^{\infty} \tilde{f}_i p_i(x)$ with $\tilde{f}_i := \frac{(f|p_i)_w}{\|p_i\|_w^2}$.

The truncated sum

$$\Pi_N^w f(x) := \sum_{i=0}^N \tilde{f}_i p_i(x)$$

is a polynomial of degree N : it is the **orthogonal projection** of f onto the finite dimensional subspace \mathbb{P}_N with respect to the scalar product $(.|.)_w$.

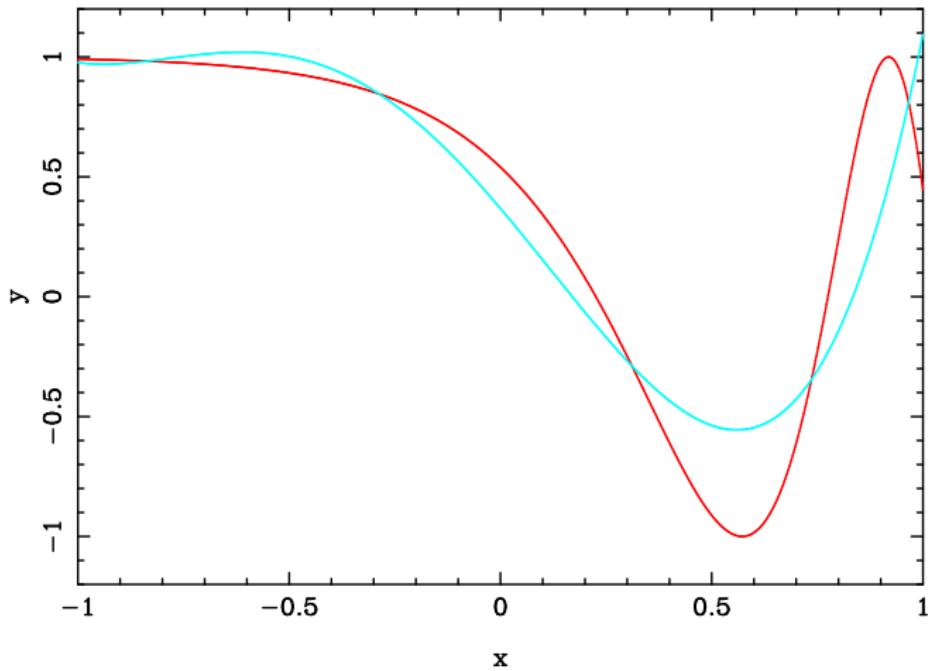
We have

$$\lim_{N \rightarrow \infty} \|f - \Pi_N^w f\|_w = 0$$

Hence $\Pi_N^w f$ can be considered as a polynomial approximation of the function f .

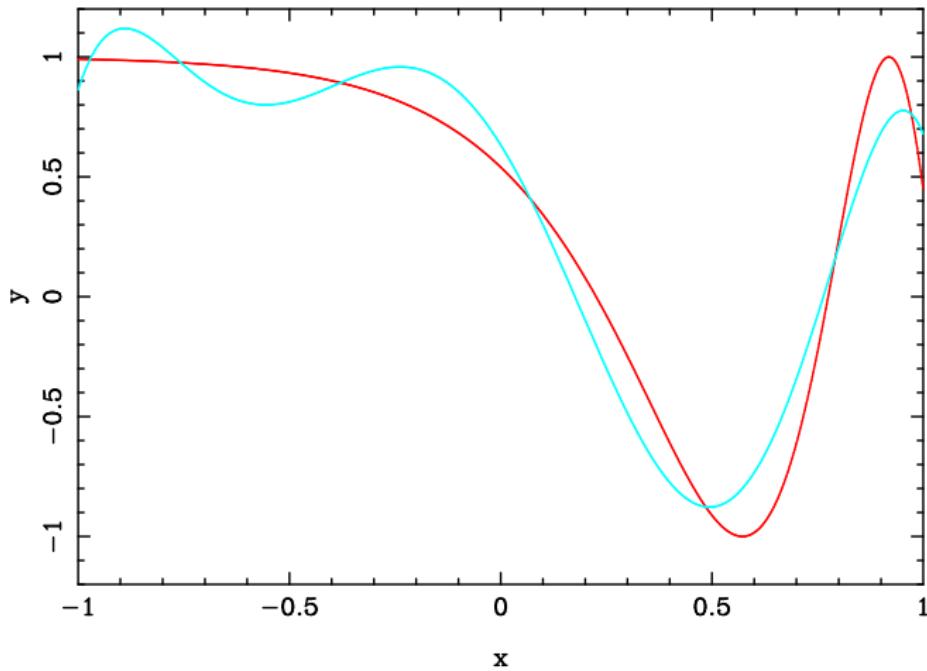
Example: Chebyshev projection of $f(x) = \cos(2 \exp(x))$

$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 4 : \|f - \Pi_4^w f\|_\infty \simeq 0.66$$



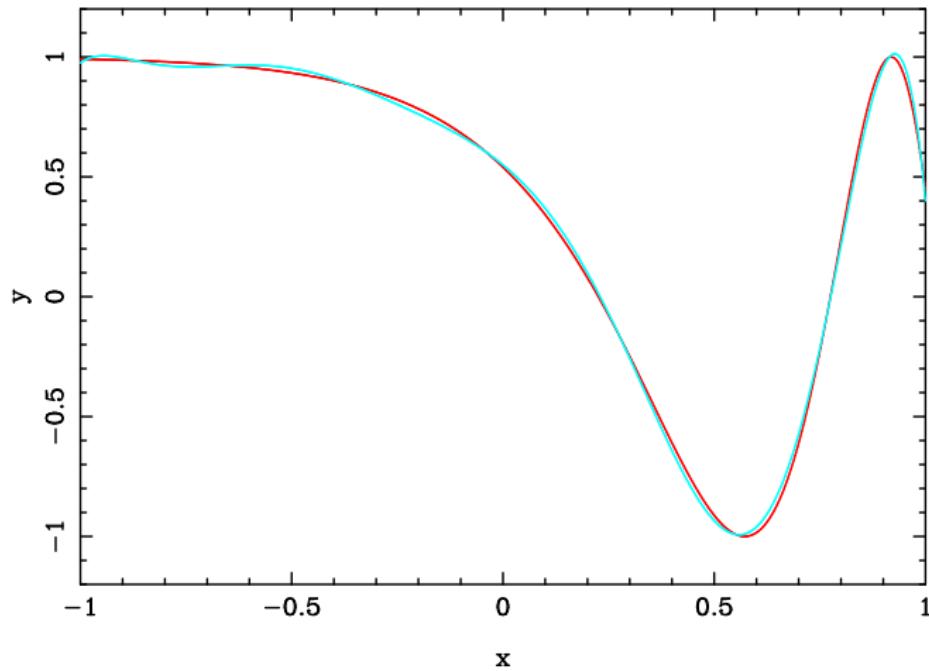
Example: Chebyshev projection of $f(x) = \cos(2 \exp(x))$

$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 6 : \|f - \Pi_6^w f\|_\infty \simeq 0.30$$



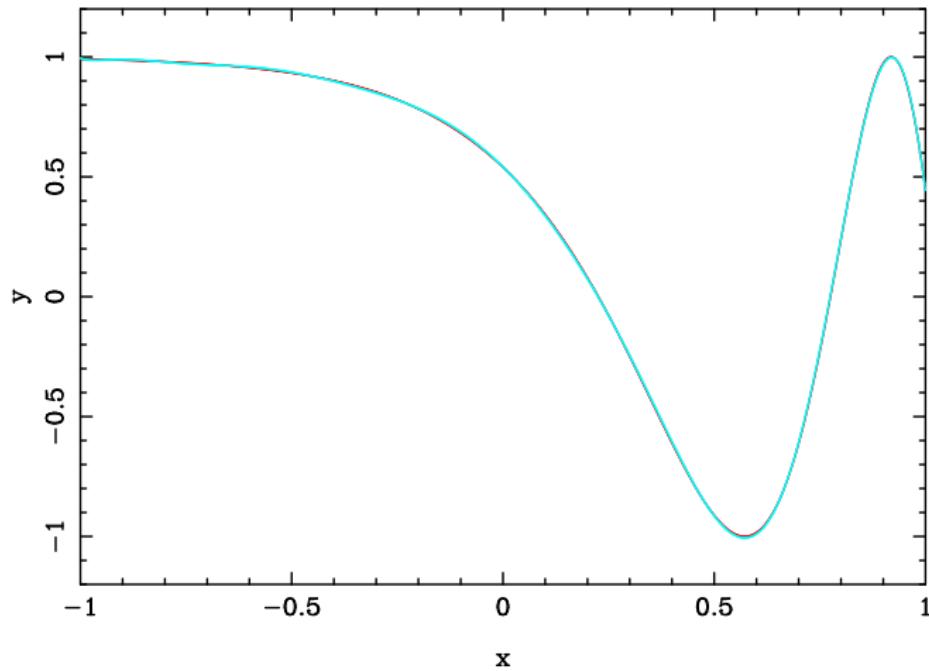
Example: Chebyshev projection of $f(x) = \cos(2 \exp(x))$

$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 8 : \|f - \Pi_8^w f\|_\infty \simeq 4.9 \cdot 10^{-2}$$



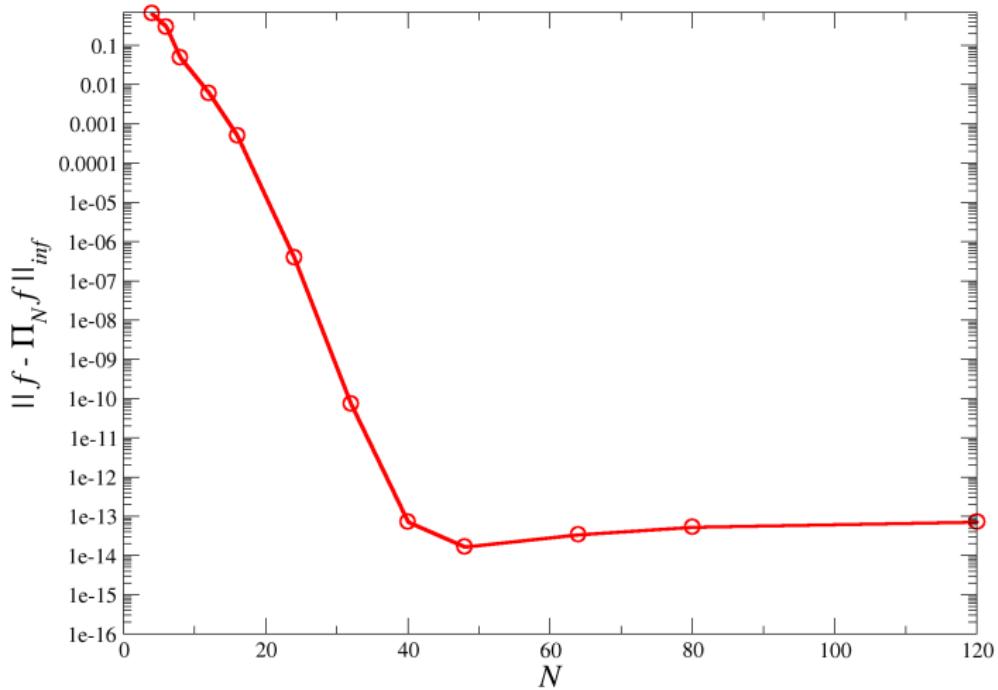
Example: Chebyshev projection of $f(x) = \cos(2 \exp(x))$

$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 12 : \|f - \Pi_{12}^w f\|_\infty \simeq 6.1 \cdot 10^{-3}$$



Example: Chebyshev projection of $f(x) = \cos(2 \exp(x))$

Variation of the projection error $\|f - \Pi_N^w f\|_\infty$ as N increases



Evaluation of the coefficients

The coefficients \tilde{f}_i of the orthogonal projection of f are given by

$$\tilde{f}_i := \frac{(f|p_i)_w}{\|p_i\|_w^2} = \frac{1}{\|p_i\|_w^2} \int_{-1}^1 f(x) p_i(x) w(x) dx \quad (2)$$

Problem: the above integral cannot be computed exactly; we must seek a numerical approximation.

Solution: Gaussian quadrature

Gaussian quadrature

Theorem (Gauss, Jacobi)

Let $(p_i)_{i \in \mathbb{N}}$ be a family of orthogonal polynomials with respect to some weight w . For $N > 0$, let $X = (x_i)_{0 \leq i \leq N}$ be the grid formed by the $N + 1$ zeros of the polynomial p_{N+1} and

$$w_i := \int_{-1}^1 \ell_i^X(x) w(x) dx$$

where ℓ_i^X is the i -th Lagrange cardinal polynomial of the grid X

◀ reminder

Then

$$\forall f \in \mathbb{P}_{2N+1}, \int_{-1}^1 f(x) w(x) dx = \sum_{i=0}^N w_i f(x_i)$$

If $f \notin \mathbb{P}_{2N+1}$, the above formula provides a good approximation of the integral.

Gauss-Lobatto quadrature

The nodes of the Gauss quadrature, being the zeros of p_{N+1} , do not encompass the boundaries -1 and 1 of the interval $[-1, 1]$. For numerical purpose, it is desirable to include these points in the boundaries.

This is possible at the price of reducing by 2 units the degree of exactness of the Gauss quadrature

Gauss-Lobatto quadrature

Theorem (Gauss-Lobatto quadrature)

Let $(p_i)_{i \in \mathbb{N}}$ be a family of orthogonal polynomials with respect to some weight w . For $N > 0$, let $X = (x_i)_{0 \leq i \leq N}$ be the grid formed by the $N + 1$ zeros of the polynomial

$$q_{N+1} = p_{N+1} + \alpha p_N + \beta p_{N-1}$$

where the coefficients α and β are such that $x_0 = -1$ and $x_N = 1$.

Let

$$w_i := \int_{-1}^1 \ell_i^X(x) w(x) dx$$

where ℓ_i^X is the i -th Lagrange cardinal polynomial of the grid X .

Then

$$\boxed{\forall f \in \mathbb{P}_{2N-1}, \int_{-1}^1 f(x) w(x) dx = \sum_{i=0}^N w_i f(x_i)}$$

Notice: $f \in \mathbb{P}_{2N-1}$ instead of $f \in \mathbb{P}_{2N+1}$ for Gauss quadrature.

Gauss-Lobatto quadrature

Remark: if the (p_i) are Jacobi polynomials, i.e. if $w(x) = (1-x)^\alpha(1+x)^\beta$, then the Gauss-Lobatto nodes which are strictly inside $(-1, 1)$, i.e. x_1, \dots, x_{N-1} , are the $N - 1$ zeros of the polynomial p'_N , or equivalently the points where the polynomial p_N is extremal.

This of course holds for Legendre and Chebyshev polynomials.

For Chebyshev polynomials, the Gauss-Lobatto nodes and weights have simple expressions:

$$x_i = -\cos \frac{\pi i}{N}, \quad 0 \leq i \leq N$$

$$w_0 = w_N = \frac{\pi}{2N}, \quad w_i = \frac{\pi}{N}, \quad 1 \leq i \leq N-1$$

Note: in the following, we consider only Gauss-Lobatto quadratures

Discrete scalar product

The Gauss-Lobatto quadrature motivates the introduction of the following scalar product:

$$\langle f|g \rangle_N = \sum_{i=0}^N w_i f(x_i) g(x_i)$$

It is called the **discrete scalar product** associated with the Gauss-Lobatto nodes $X = (x_i)_{0 \leq i \leq N}$

Setting $\gamma_i := \langle p_i | p_i \rangle_N$, the **discrete coefficients** associated with a function f are given by

$$\hat{f}_i := \frac{1}{\gamma_i} \langle f | p_i \rangle_N, \quad 0 \leq i \leq N$$

which can be seen as approximate values of the coefficients \tilde{f}_i provided by the Gauss-Lobatto quadrature [cf. Eq. (2)]

Discrete coefficients and interpolating polynomial

Let $I_N^{\text{GL}} f$ be the interpolant of f at the Gauss-Lobatto nodes $X = (x_i)_{0 \leq i \leq N}$. Being a polynomial of degree N , it is expandable as

$$I_N^{\text{GL}} f(x) = \sum_{i=0}^N a_i p_i(x)$$

Then, since $I_N^{\text{GL}} f(x_j) = f(x_j)$,

$$\hat{f}_i = \frac{1}{\gamma_i} \langle f | p_i \rangle_N = \frac{1}{\gamma_i} \langle I_N^{\text{GL}} f | p_i \rangle_N = \frac{1}{\gamma_i} \sum_{j=0}^N a_j \langle p_j | p_i \rangle_N$$

Now, if $j = i$, $\langle p_j | p_i \rangle_N = \gamma_i$ by definition. If $j \neq i$, $p_j p_i \in \mathbb{P}_{2N-1}$ so that the Gauss-Lobatto formula holds and gives $\langle p_j | p_i \rangle_N = (p_j | p_i)_w = 0$. Thus we conclude that $\langle p_j | p_i \rangle_N = \gamma_i \delta_{ij}$ so that the above equation yields $\hat{f}_i = a_i$, i.e. **the discrete coefficients are nothing but the coefficients of the expansion of the interpolant at the Gauss-Lobato nodes**

Spectral representation of a function

In a spectral method, the numerical representation of a function f is through its interpolant at the Gauss-Lobatto nodes:

$$I_N^{\text{GL}} f(x) = \sum_{i=0}^N \hat{f}_i p_i(x)$$

The discrete coefficients \hat{f}_i are computed as

$$\hat{f}_i = \frac{1}{\gamma_i} \sum_{j=0}^N w_j f(x_j) p_i(x_j)$$

$I_N^{\text{GL}} f(x)$ is an approximation of the truncated series $\Pi_N^w f(x) = \sum_{i=0}^N \tilde{f}_i p_i(x)$,

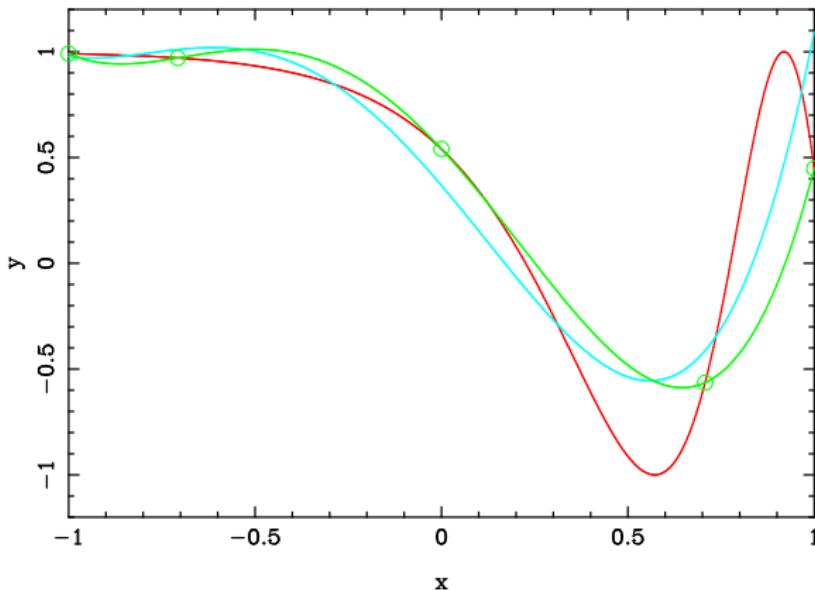
which is the orthogonal projection of f onto the polynomial space \mathbb{P}_N .

$\Pi_N^w f$ should be the true spectral representation of f , but in general it is not computable exactly.

The difference between $I_N^{\text{GL}} f$ and $\Pi_N^w f$ is called the **aliasing error**

Example: aliasing error for $f(x) = \cos(2 \exp(x))$

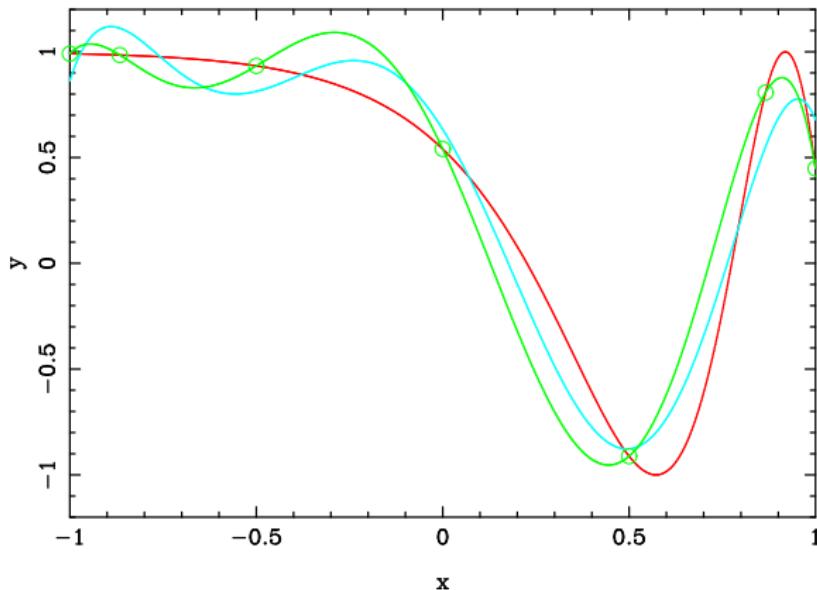
$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 4$$



red: f ; blue: $\Pi_N^w f$; green: $I_N^{\text{GL}} f$

Example: aliasing error for $f(x) = \cos(2 \exp(x))$

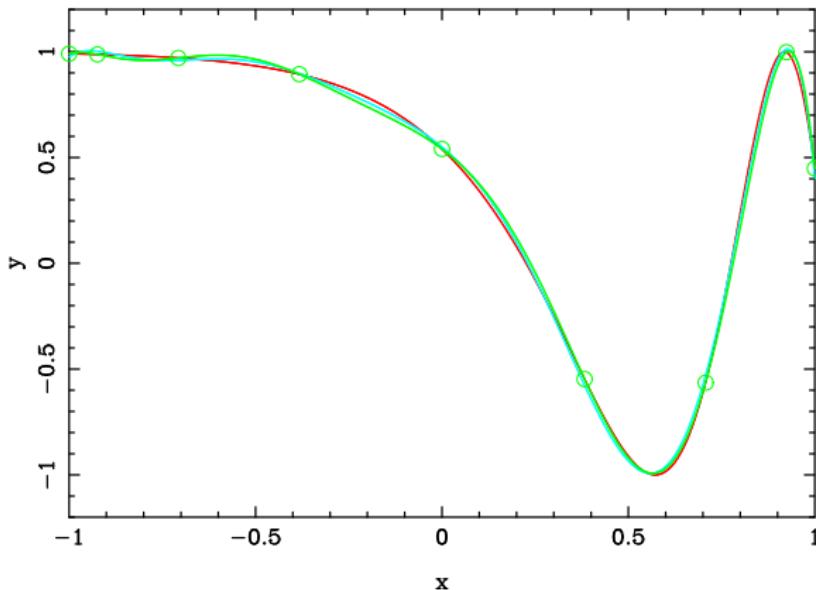
$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 6$$



red: f ; blue: $\Pi_N^w f$; green: $I_N^{\text{GL}} f$

Example: aliasing error for $f(x) = \cos(2 \exp(x))$

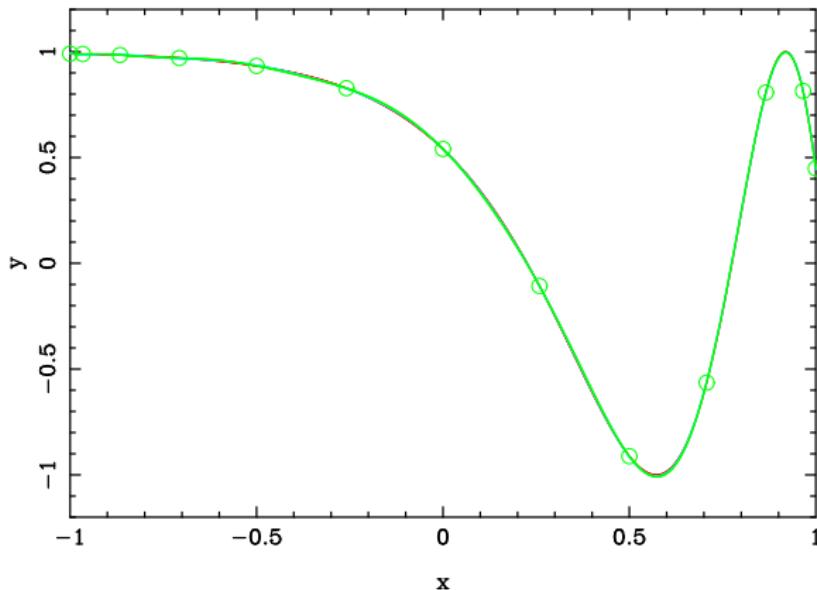
$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 8$$



red: f ; blue: $\Pi_N^w f$; green: $I_N^{\text{GL}} f$

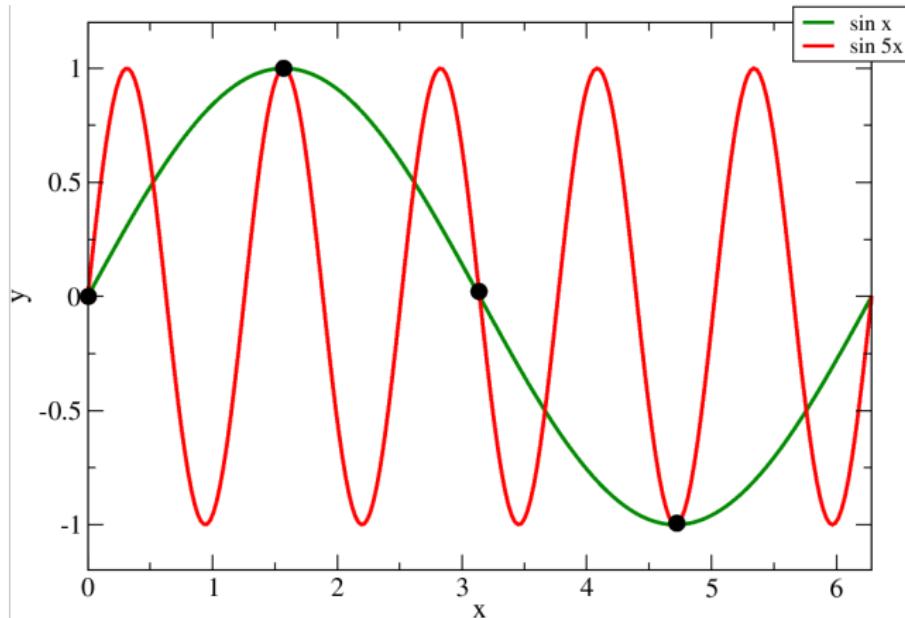
Example: aliasing error for $f(x) = \cos(2 \exp(x))$

$$f(x) = \cos(2 \exp(x)) \quad w(x) = (1 - x^2)^{-1/2} \quad N = 12$$



red: f ; blue: $\Pi_N^w f$; green: $I_N^{\text{GL}} f$

Aliasing error = contamination by high frequencies



Aliasing of a $\sin(x)$ wave by a $\sin(5x)$ wave on a 4-points grid

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- 1 Introduction
- 2 Interpolation on an arbitrary grid
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- 4 Convergence of the spectral expansions
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Sobolev norm

Let us consider a function $f \in C^m([-1, 1])$, with $m \geq 0$.

The **Sobolev norm** of f with respect to some weight function w is

$$\|f\|_{H_w^m} := \left(\sum_{k=0}^m \|f^{(k)}\|_w^2 \right)^{1/2}$$

Convergence rates for $f \in C^m([-1, 1])$

Chebyshev expansions:

- truncation error :

$$\|f - \Pi_N^w f\|_w \leq \frac{C_1}{N^m} \|f\|_{H_w^m} \text{ and } \|f - \Pi_N^w f\|_\infty \leq \frac{C_2(1 + \ln N)}{N^m} \sum_{k=0}^m \|f^{(k)}\|_\infty$$

- interpolation error :

$$\|f - I_N^{\text{GL}} f\|_w \leq \frac{C_3}{N^m} \|f\|_{H_w^m} \text{ and } \|f - I_N^{\text{GL}} f\|_\infty \leq \frac{C_4}{N^{m-1/2}} \|f\|_{H_w^m}$$

Legendre expansions:

- truncation error :

$$\|f - \Pi_N^w f\|_w \leq \frac{C_1}{N^m} \|f\|_{H_w^m} \text{ and } \|f - \Pi_N^w f\|_\infty \leq \frac{C_2}{N^{m-1/2}} V(f^{(m)})$$

- interpolation error :

$$\|f - I_N^{\text{GL}} f\|_w \leq \frac{C_3}{N^{m-1/2}} \|f\|_{H_w^m}$$

Evanescence error for smooth functions

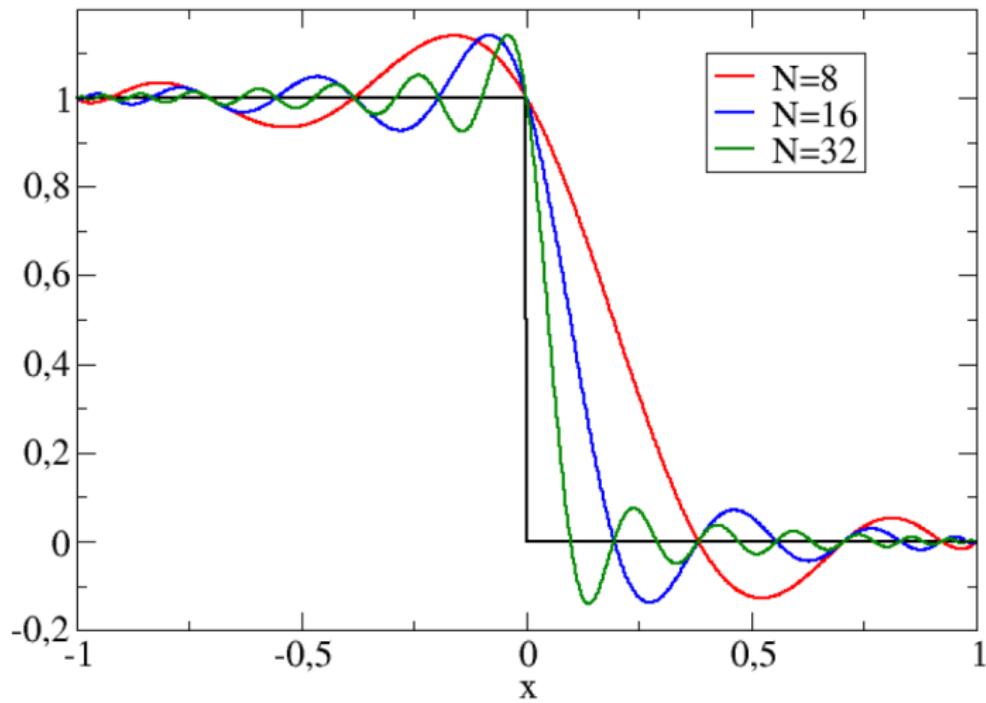
If $f \in C^\infty([-1, 1])$, the error of the spectral expansions $\Pi_N^w f$ or $I_N^{GL} f$ decays more rapidly than any power of N .

In practice: **exponential decay** ◀ example

This error is called **evanescent**.

For non-smooth functions: Gibbs phenomenon

Extreme case: f discontinuous



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